

## Signal flow graphs

- ▶ Can be expressed either in time domain (delay  $T$ ) or frequency domain (delay  $z^{-1}$ )
- ▶ No delay free loops to be sequentially computable
- ▶ Fully specified: expressed in the operations that will be used for implementation
- ▶ Precedence graph: determines which operations are sequential and which can be performed concurrently
- ▶ Operations in computable order: difference equations can be derived based on the precedence graph to obtain a sequential "program"

## Application Specific Integrated Circuits for Digital Signal Processing Lecture 12

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## Today's topic

### Algorithm properties

- ▶ Parallelism
    - ▶ Recursive algorithms
- $$T_{\min} = \max_i \frac{\sum_i T_{L,i}}{N_i}$$
- ▶ Non-recursive: unlimited
  - ▶ Critical path: longest time directed path,  $T_{cp}$
- ### Operations
- ▶ Latency: delay from input to corresponding output  $\rightarrow$  scheduling,  $T_{\min}$ ,  $T_{cp}$
  - ▶ Execution time: delay from input (output) to next input (output)  $\rightarrow$  resource usage
- ### Finite word-length effects
- ▶ Scaling: avoid overflows while utilizing the complete numerical range
  - ▶ Data quantization: round-off noise, possibly stability issues for recursive algorithms
  - ▶ Coefficient quantization: static deviation

- ▶ Course review

## Algorithm transformations

- ▶ Pipelining: introduce delay elements  $\rightarrow$  change  $T_{cp}$
- ▶ Retiming: move delay elements  $\rightarrow$  change  $T_{cp}$
- ▶ Reordering and/or merging operations  $\rightarrow$  change  $T_{cp}$  and  $T_{min}$
- ▶ Lookahead pipelining and block processing  $\rightarrow$  change  $T_{min}$  (and  $T_{cp}$ ), numerical issues
- ▶ Unfolding: express algorithm over several sample periods  $\rightarrow$  does NOT change  $T_{min}$  (but may allow transformations)

## Scheduling

- ▶ When to execute which process (variable storage is also a process)
- ▶ Will determine the resource requirements
- ▶ Cyclic scheduling
  - ▶ For iterative processing, there will always be a next iteration
  - ▶ Outputs does not have to be explicitly "flushed out"
  - ▶ Can schedule on a cylinder  $\rightarrow$  no explicit start or stop time
  - ▶ Moving an operation across the scheduling boundary equivalent to pipelining/retiming
- ▶ Scheduling period
  - ▶ Integer multiple of  $T_{sample}$
  - ▶ Integer multiple of measure time (commonly clock cycles)
  - ▶ At least as long as maximum execution time (including variable storage)

## Resource allocation and assignment

- ▶ Resource allocation: determining the number of resources
- ▶ Resource assignment: determining which process/operation is using which resource
- ▶ Often performed concurrently
  - ▶ Clique partitioning: inclusion or exclusion graph, NP-hard
  - ▶ Left-edge algorithm: fast heuristic
- ▶ Resource assignment will determine communication complexity
- ▶ Memory/storage
  - ▶ Number of concurrent accesses determines the number of memories/memory ports
  - ▶ Number of concurrent variables to be stored determines the number of cells

## Architecture

- ▶ Four separate parts
    - ▶ Processing elements: performs the computations
    - ▶ Memory: store information
    - ▶ Communication: connects memory and PEs
    - ▶ Control: determines what and when things happen
  - ▶ Shared memory architecture
    - ▶  $N$  PE and  $M$  memories, each PE has exclusive access during one cycle
    - ▶ Memory bottleneck
- $$T_{exe,PE} \geq 2NT_{exe,mem}$$
- ▶ Increase  $T_{exe,PE}$ : use larger operations  $\rightarrow$  less flexible schedule, fewer memory accesses
  - ▶ Decrease  $T_{exe,mem}$ : interleave memories  $\rightarrow$  constraints on schedule
  - ▶ Decrease memory accesses: broadcasting, cache memories  $\rightarrow$  constraints on schedule

## Arithmetic

- ▶ Fixed-point representations: main problem is to represent negative numbers
- ▶ Floating-point representations: higher dynamic range, but lower resolution, more complex operations
- ▶ Adders: carry propagation
- ▶ Multiplication: handling sign-bits
- ▶ Distributed arithmetic: efficient way to compute inner products
- ▶ Constant multiplication: avoid unnecessary computations, share computations

## Course description style

- ▶ General methods and principles used for designing and implementing application specific integrated circuits are presented
- ▶ The emphasis is put on a systematic design methodology going from system specification down to complete integrated circuit
- ▶ After completing the course you are expected to be able to:
  - ▶ Analyze computational properties of signal processing algorithms
  - ▶ Determine resource requirements for implementation of signal processing algorithms
  - ▶ Determine the suitability of different types of architectures for implementation of signal processing algorithms
  - ▶ Synthesize nearly optimal architectures
  - ▶ Realize and evaluate different types of processing elements

## Course description style

- ▶ As parts of the course the student is expected to be able to:
  - ▶ Determine limits on the obtainable data rate
  - ▶ Determine a computational order of operations
  - ▶ Compute the latency and execution time of operations and determine the impact of these properties
  - ▶ Determine computational graphs and optimize schedules
  - ▶ Define and analyze different types of optimality in parts of the design flow
  - ▶ Determine and optimize resource allocation and resource assignment
  - ▶ Synthesize and analyze different architectures for signal processing systems
  - ▶ Realize different types of processing elements using different arithmetic styles