

Lecture 8

- A memory controller
- Lab4 – a special instruction

Guest lecture

Time: Friday 5/12 1515-1600.

Place: Nollstället

Image Processing on FPGAs.

Johan Pettersson, Sick IVP

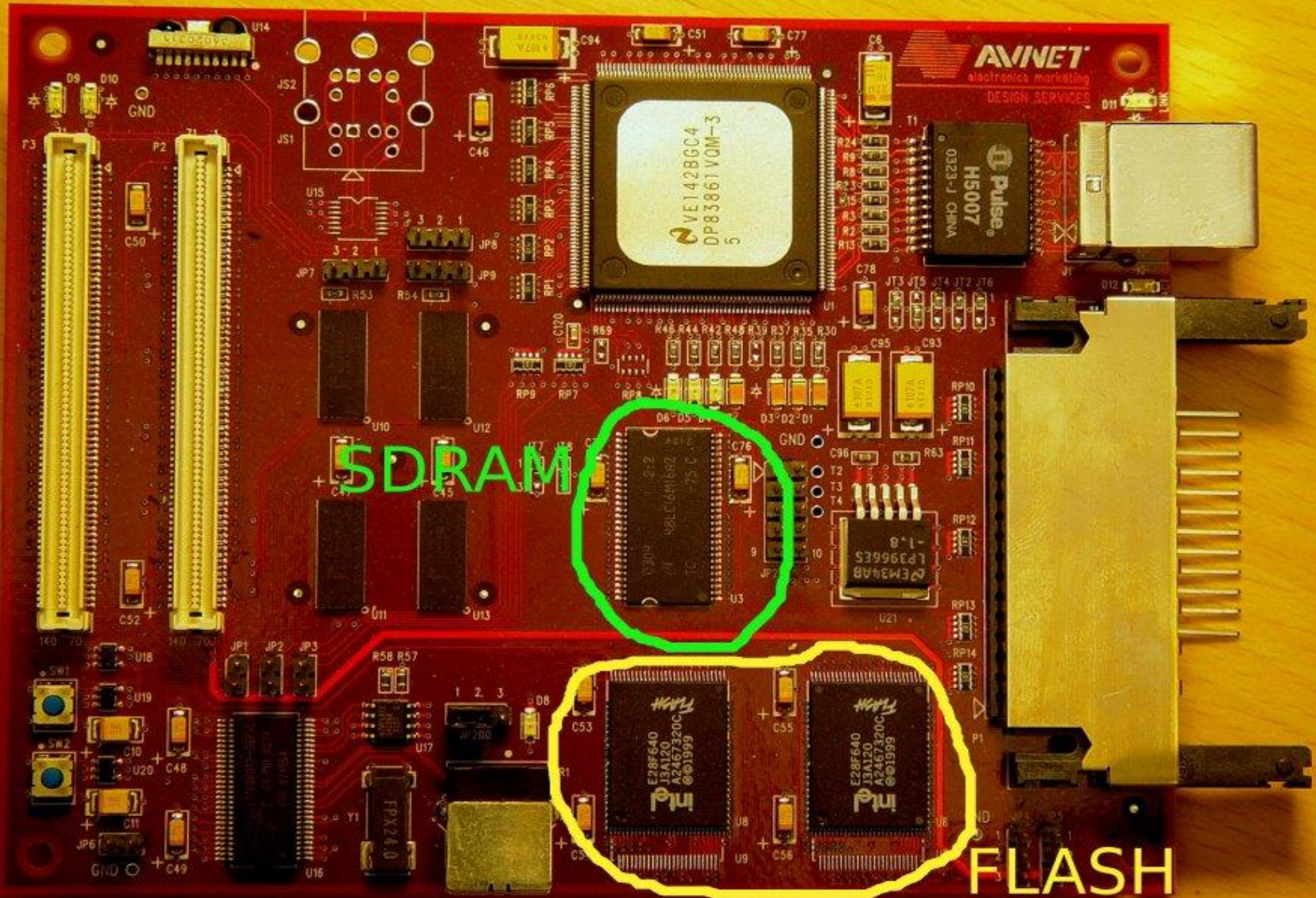
PKMC

Wishbone bus

Memory bus



- 3 controllers in one
- no registers
 - SDRAM on 0x0,
 - SRAM on 0x20000000 or 0xc0000000
 - Flash on 0xf0000000



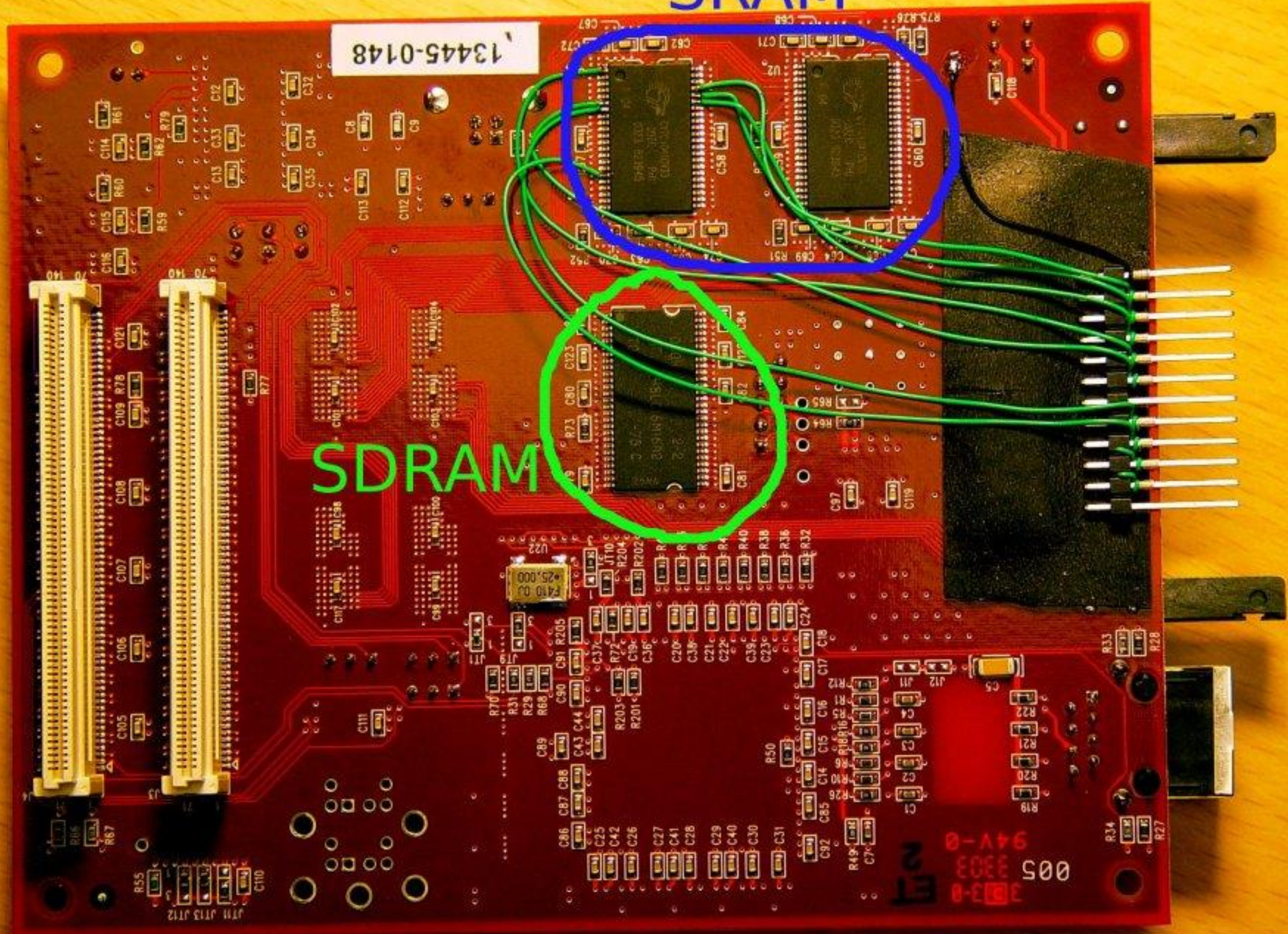
SDRAM

FLASH

SRAM

SDRAM

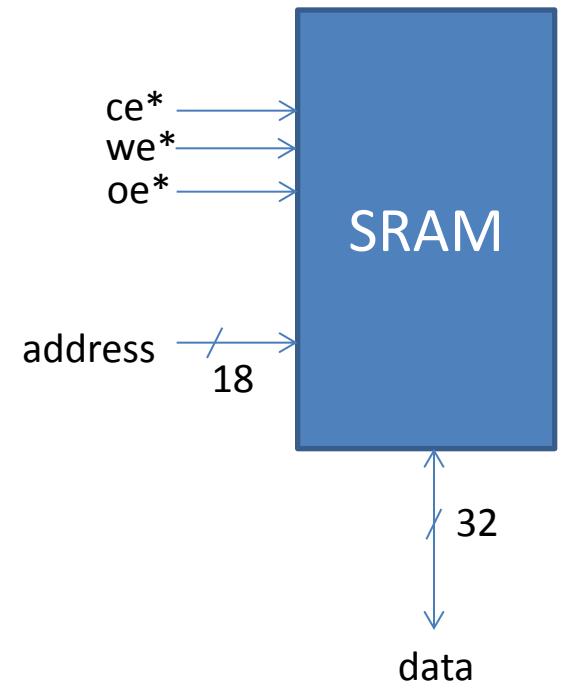
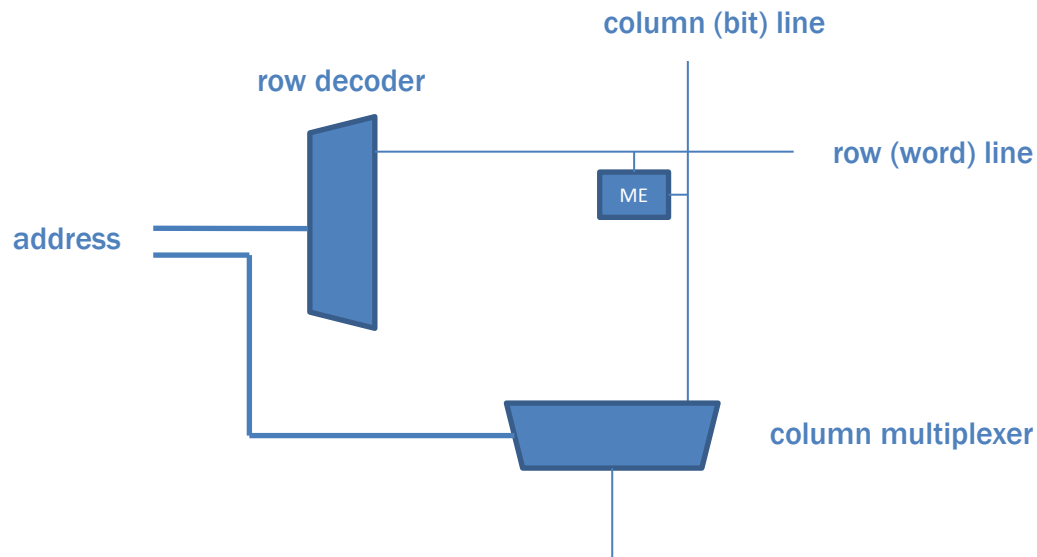
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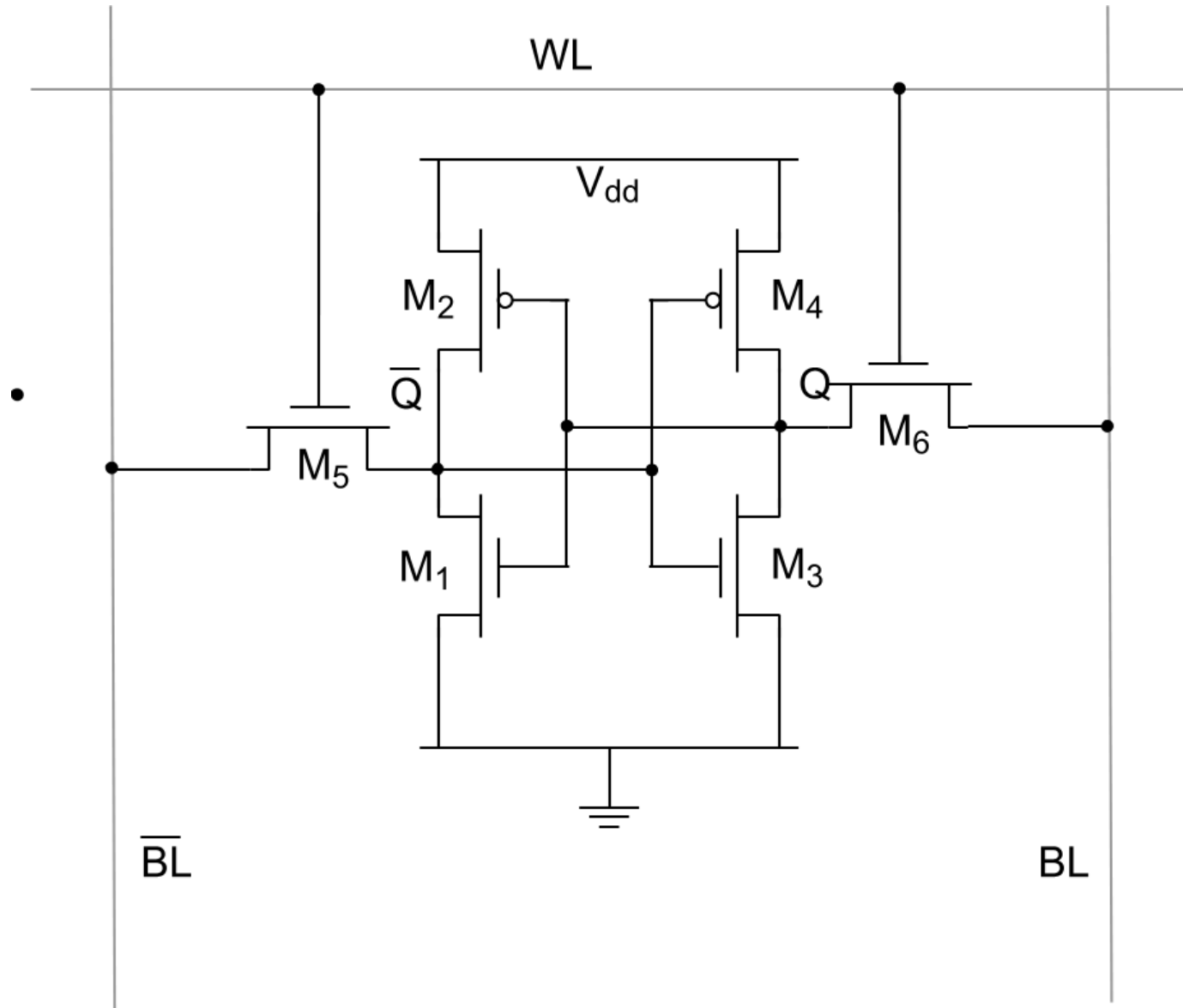
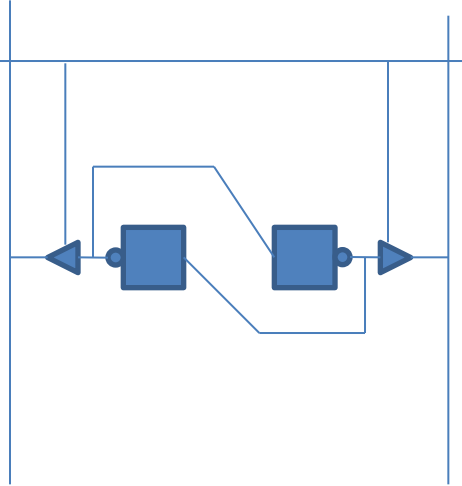
SRAM

Static RAM

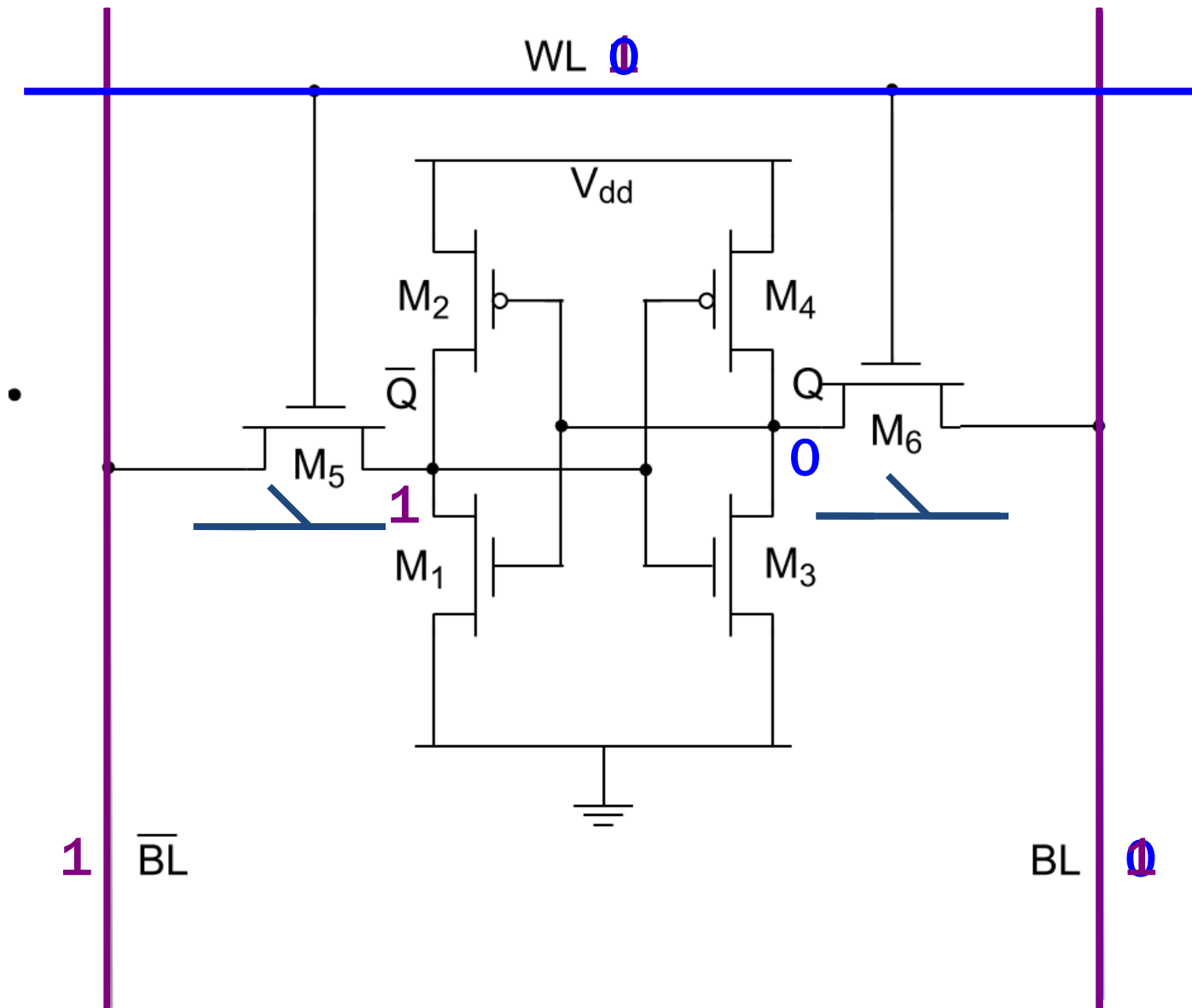
- Asynchronous device
- Memory element: Latch
- $2 \times (256k \times 16) = 1MB$



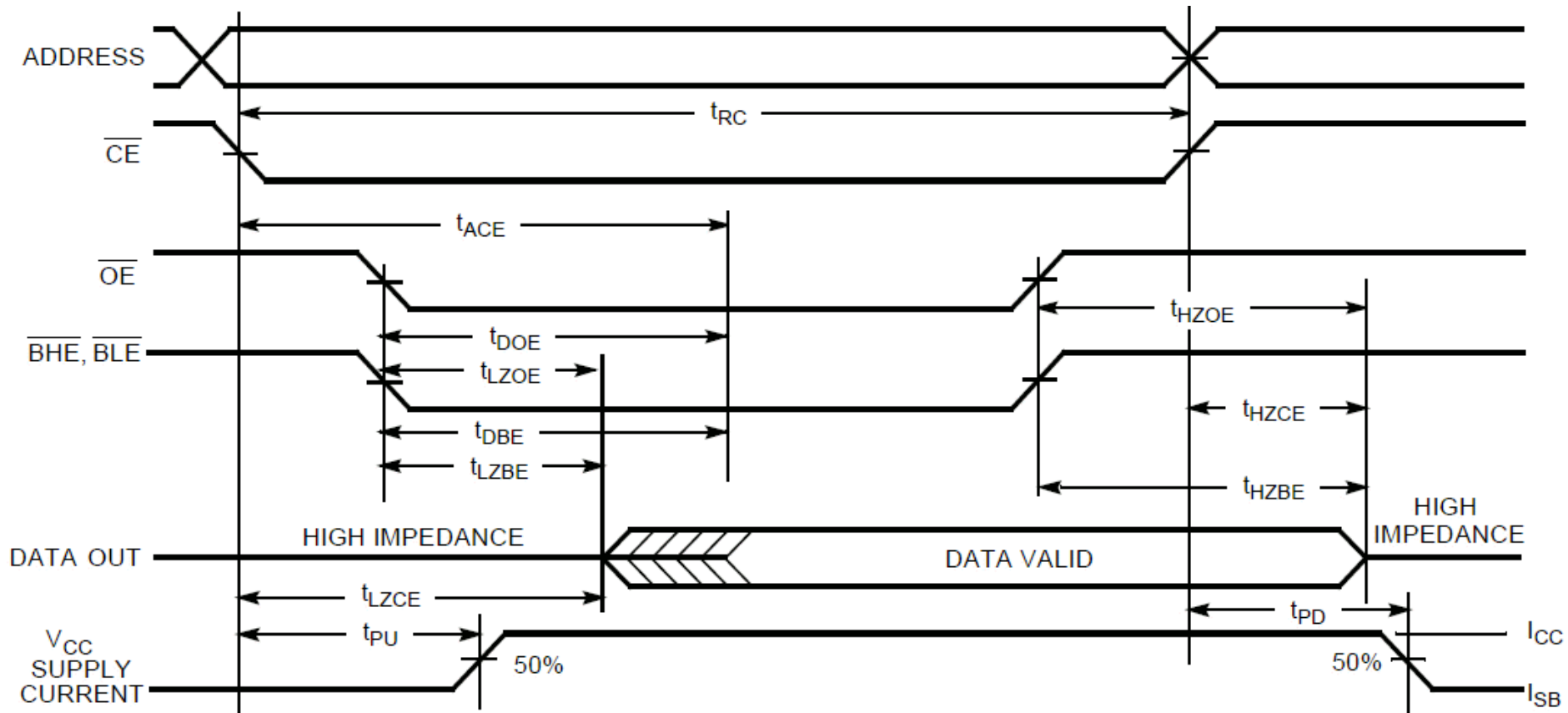
SRAM - 6 transistor cell



SRAM - Read

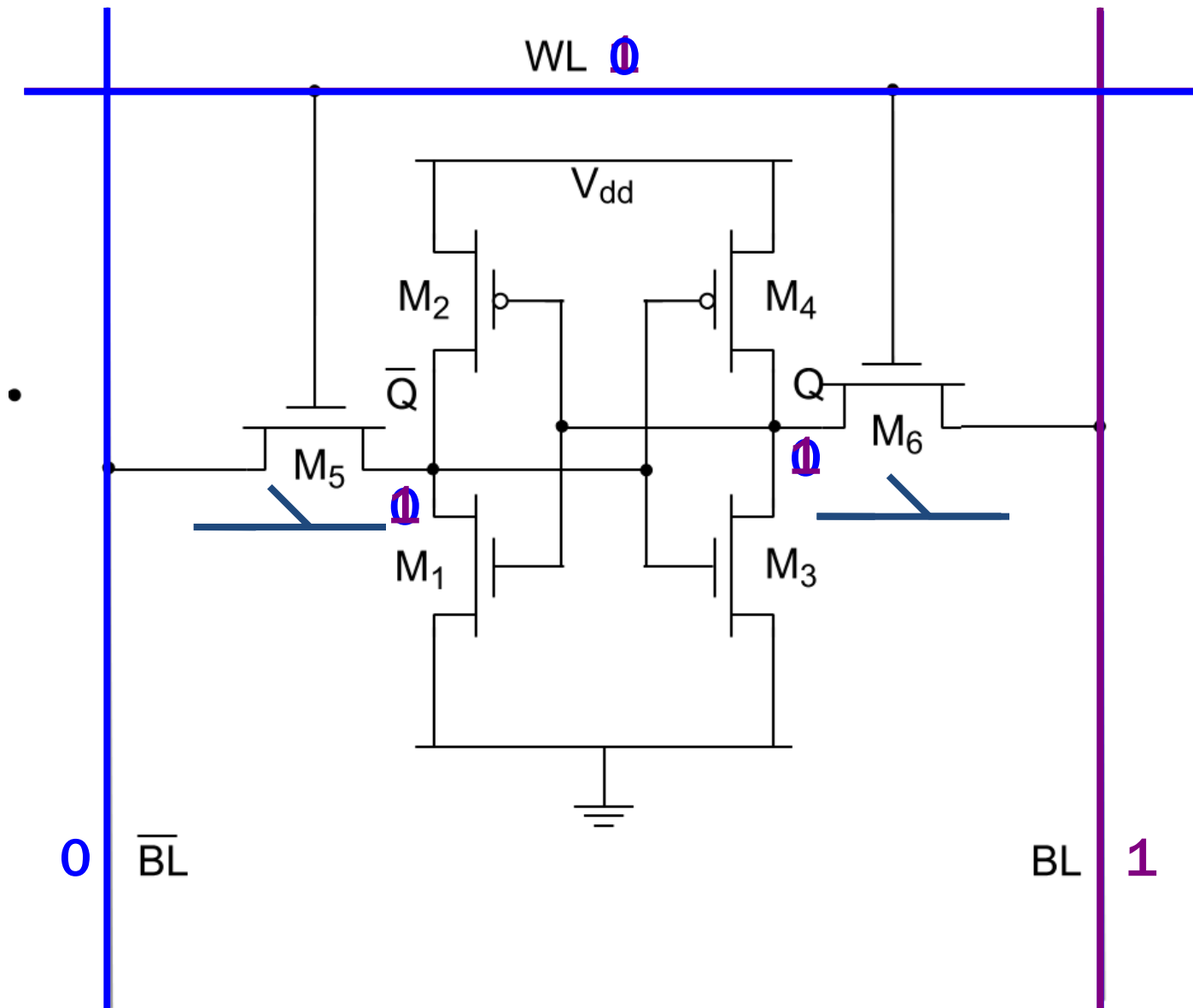


SRAM - Read

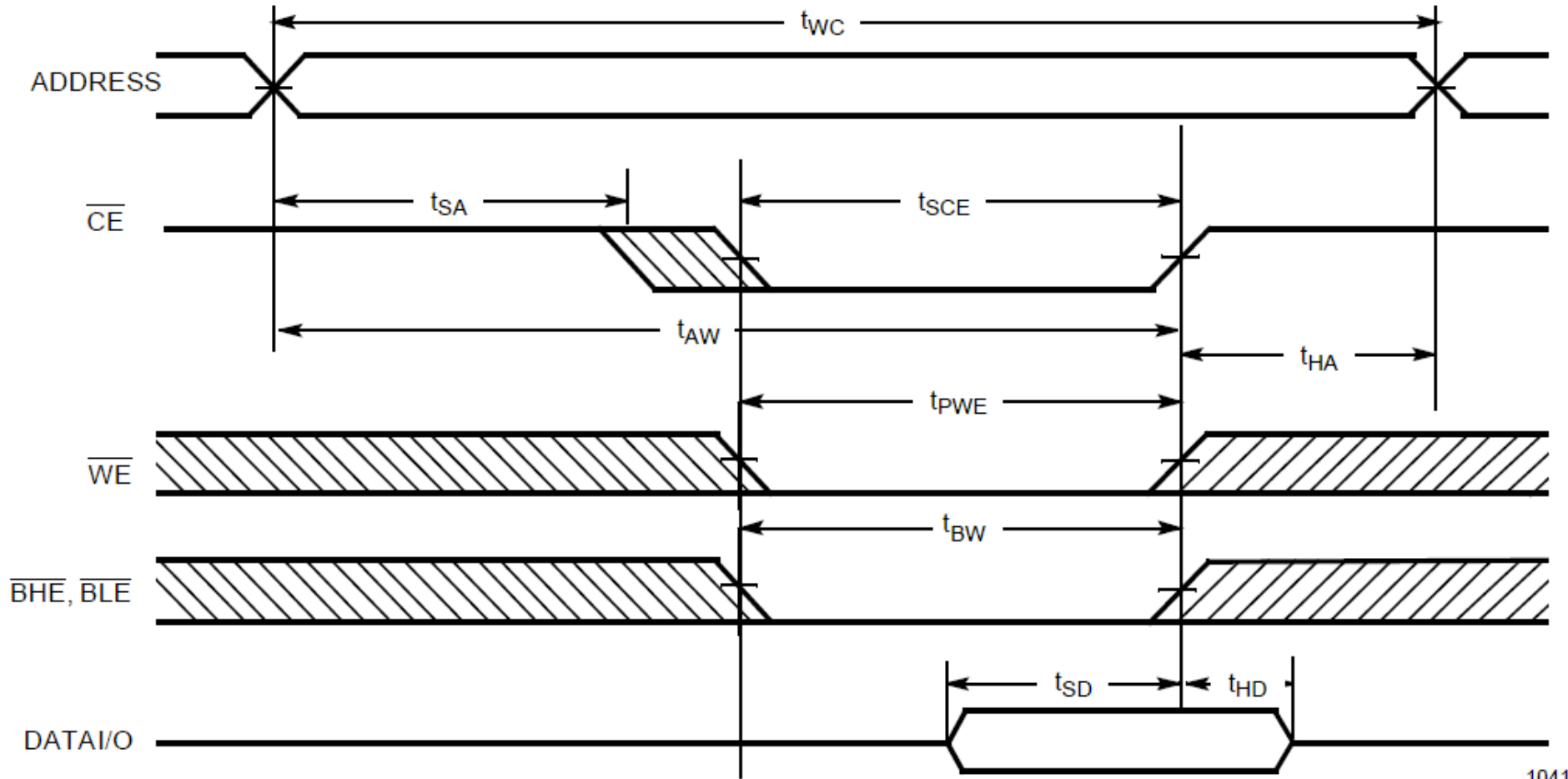


1041V33-7

SRAM - Write

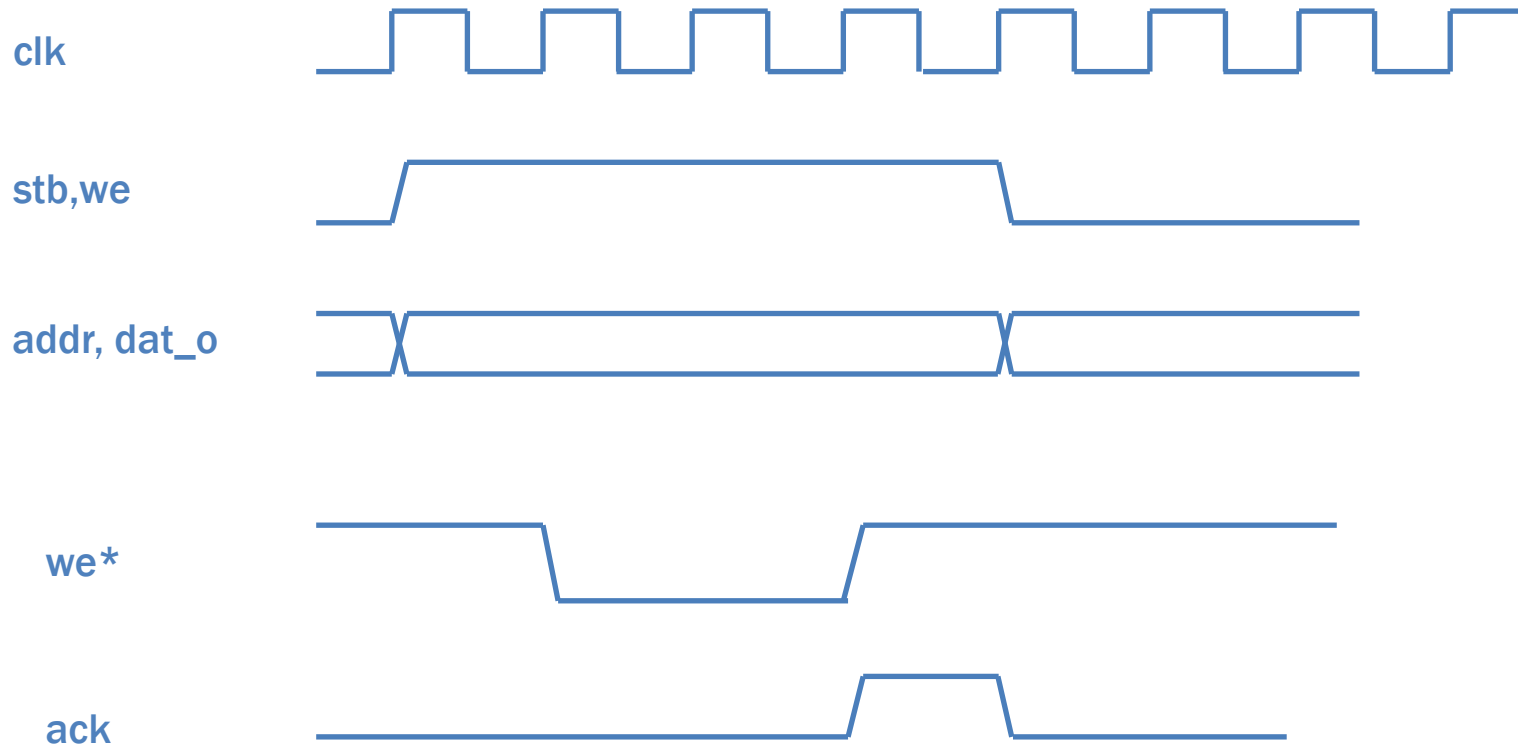


SRAM - Write



1041

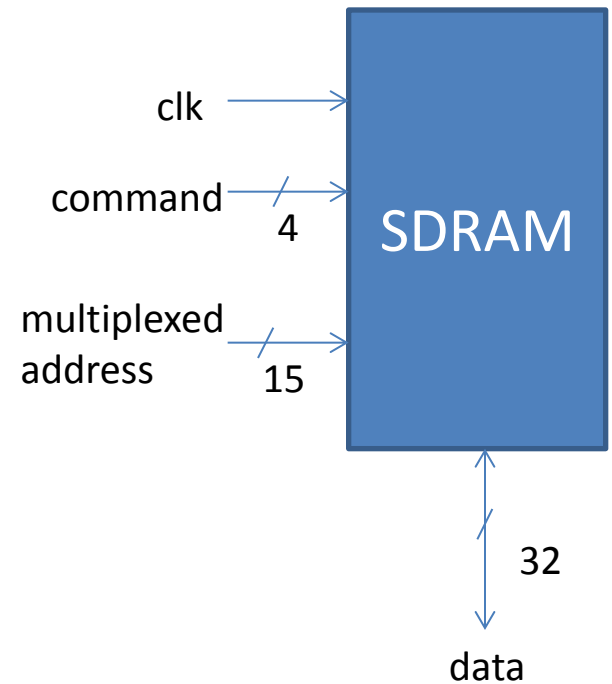
Some signals from PKMC



SDRAM

Synchronous Dynamic RAM

- Clocked device
- Memory element: Capacitance
- Needs periodic refreshing
- Pipelined operation
 - Needs a series of commands
 - No loss using muxed address
- Burst oriented
 - One address many data
- $2 \times (16M \times 16) = 64MB$



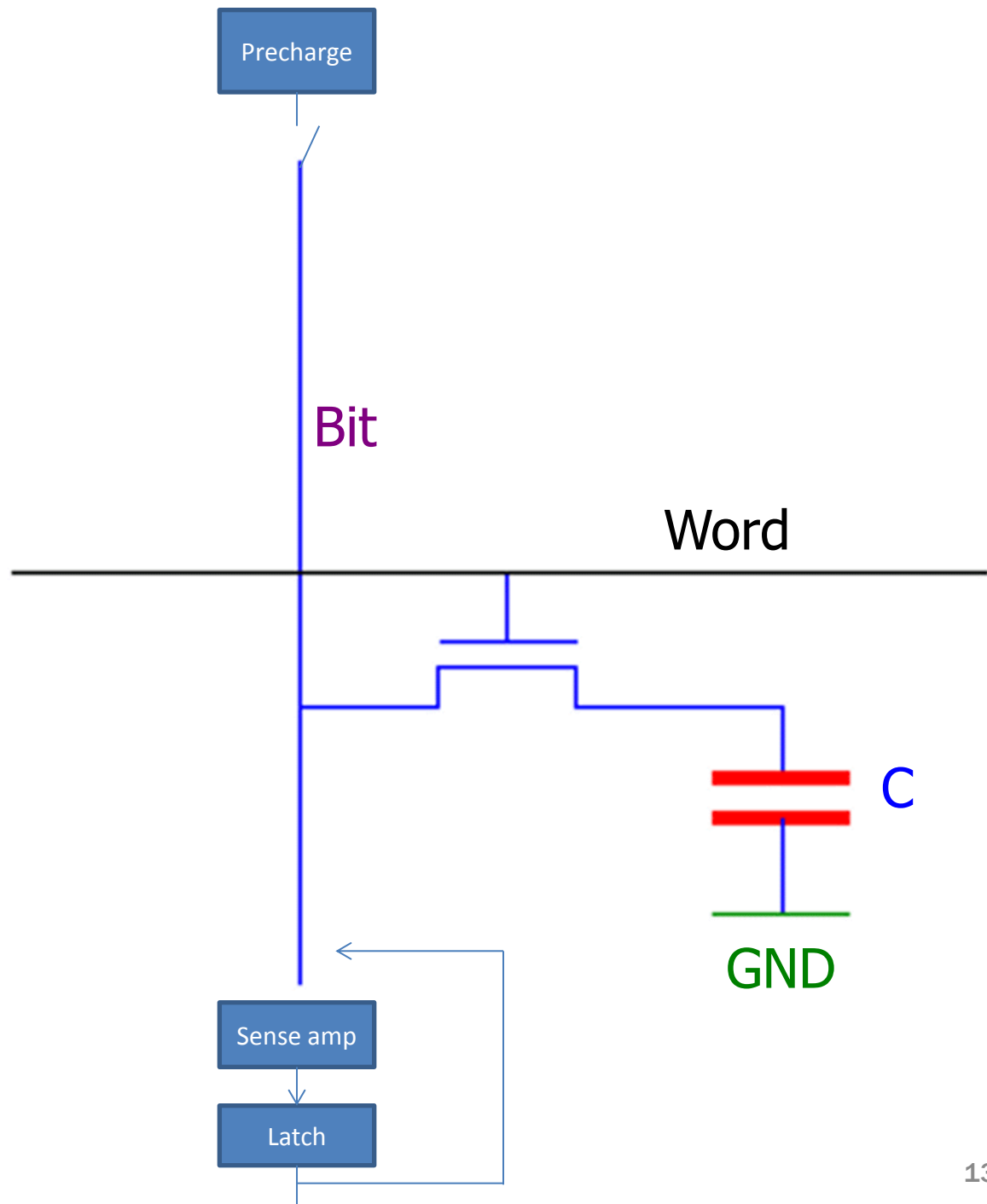
SDRAM

Sketchy read cycle

1. Precharge bit line to $V/2$
2. Let bit line float
3. Connect sense amp to bit line
4. Connect C to bit line
5. Hold value in latch
6. Write value back to C

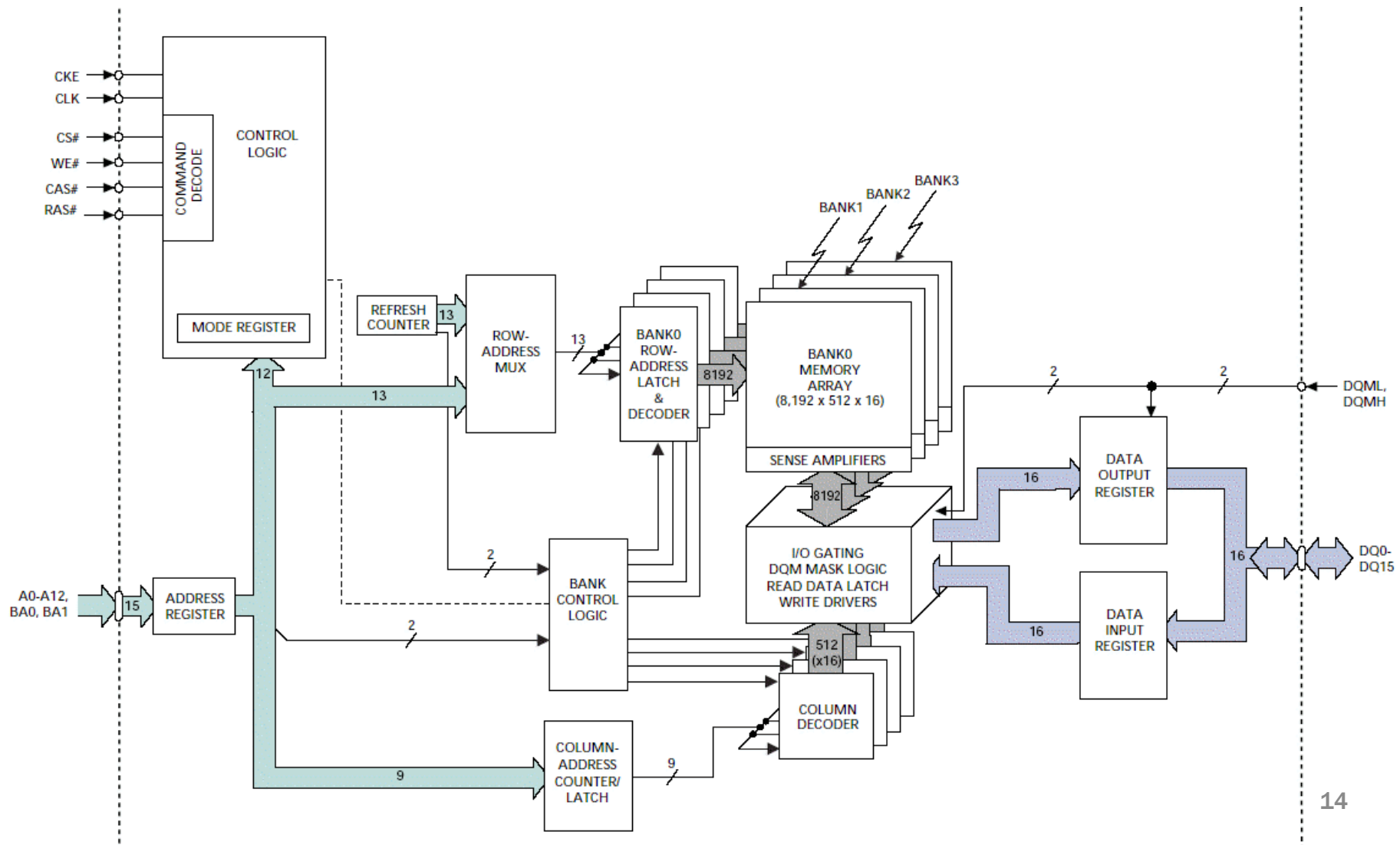
Refresh cycle

Dummy read cycle for a whole row

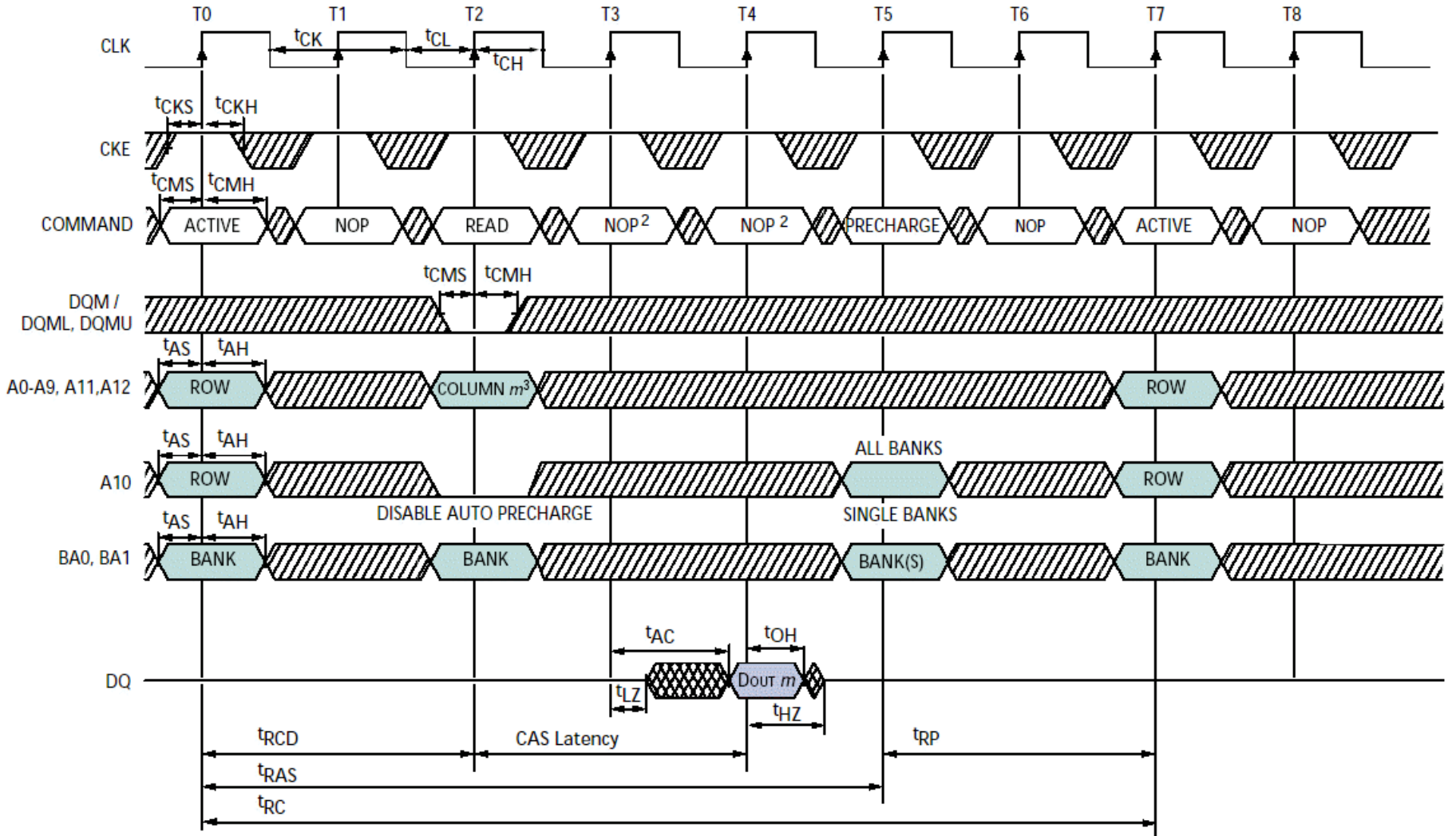




SDRAM Architecture

32 MB = 16 MW per chip, address bits = 24 = 13+11
burst oriented

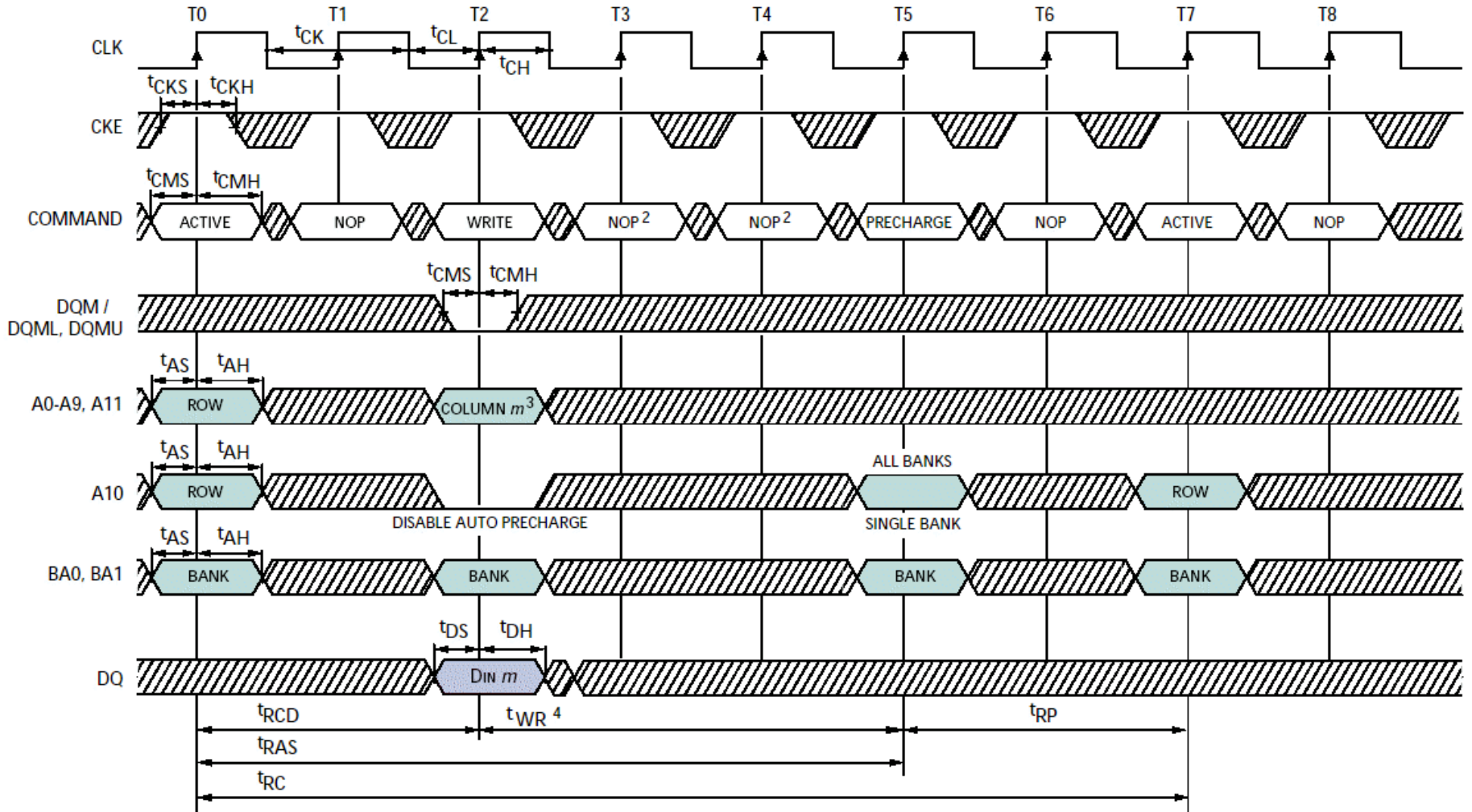


SDRAM Read



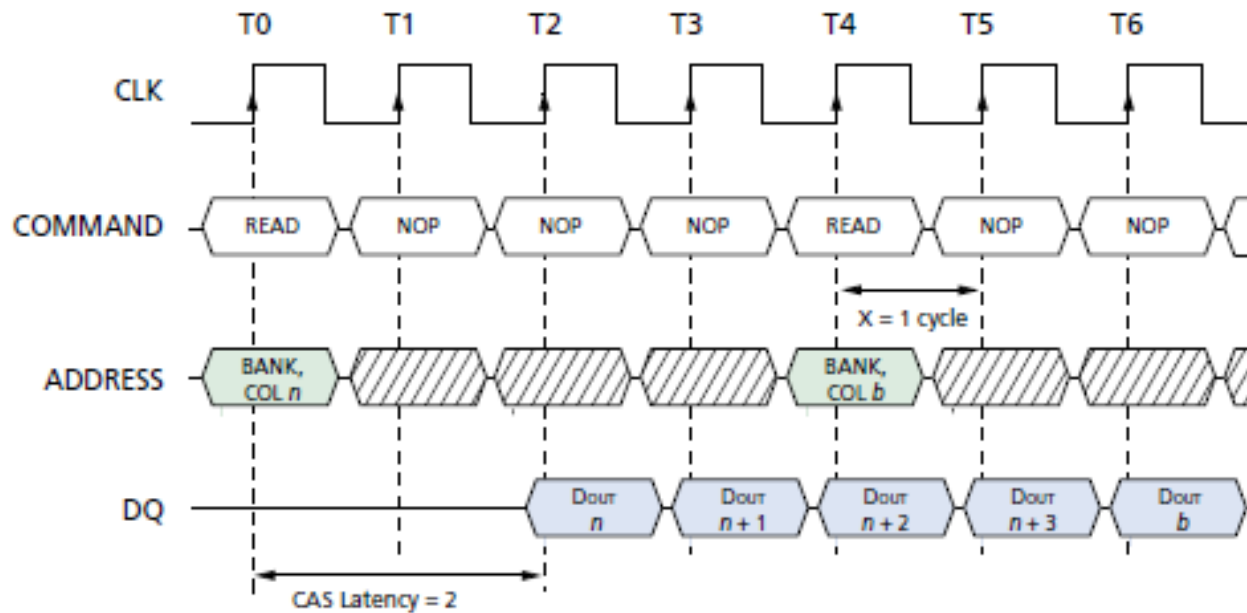
 DON'T CARE
 15
 UNDEFINED

SDRAM Write



Consecutive read bursts

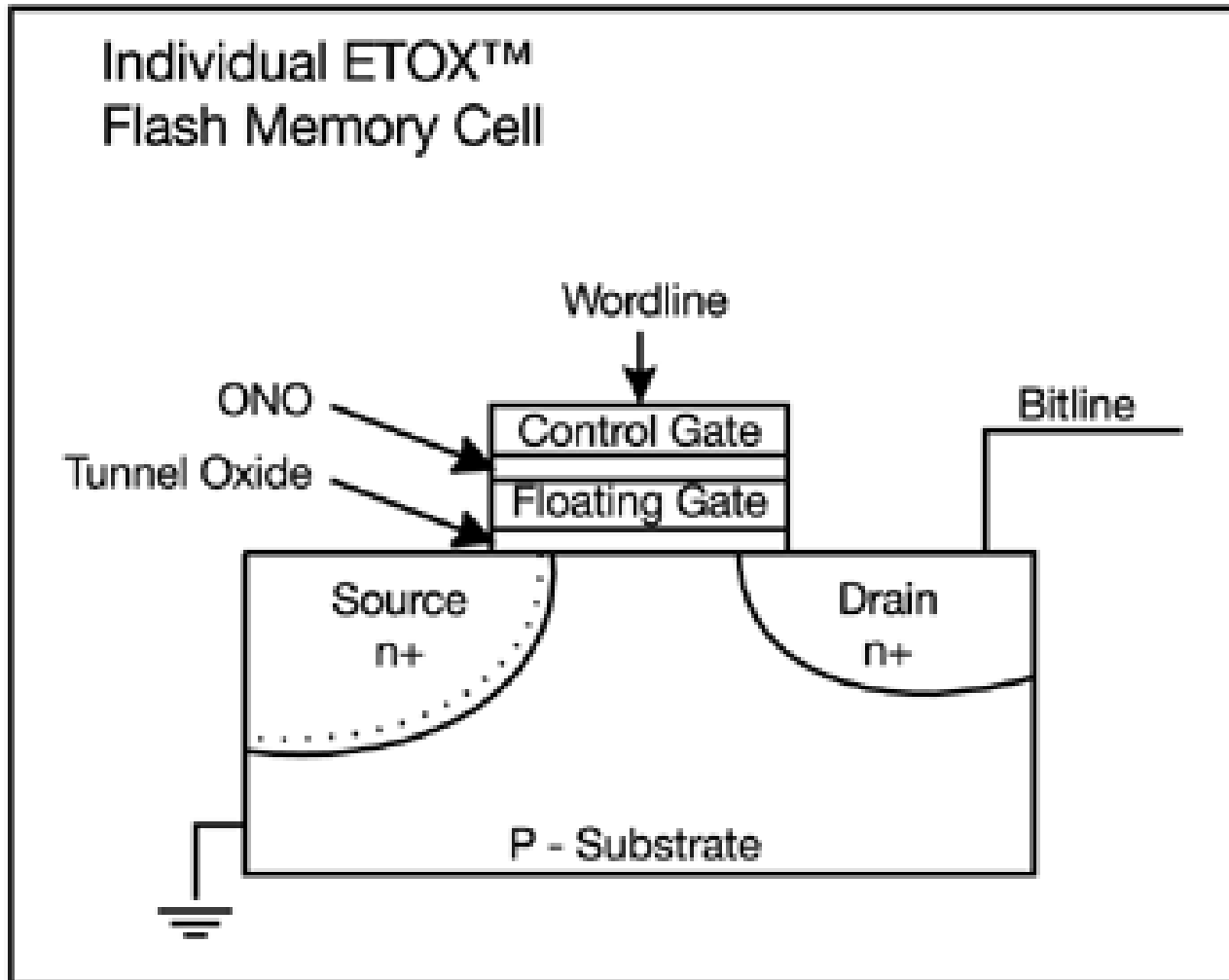
Mode register must be programmed



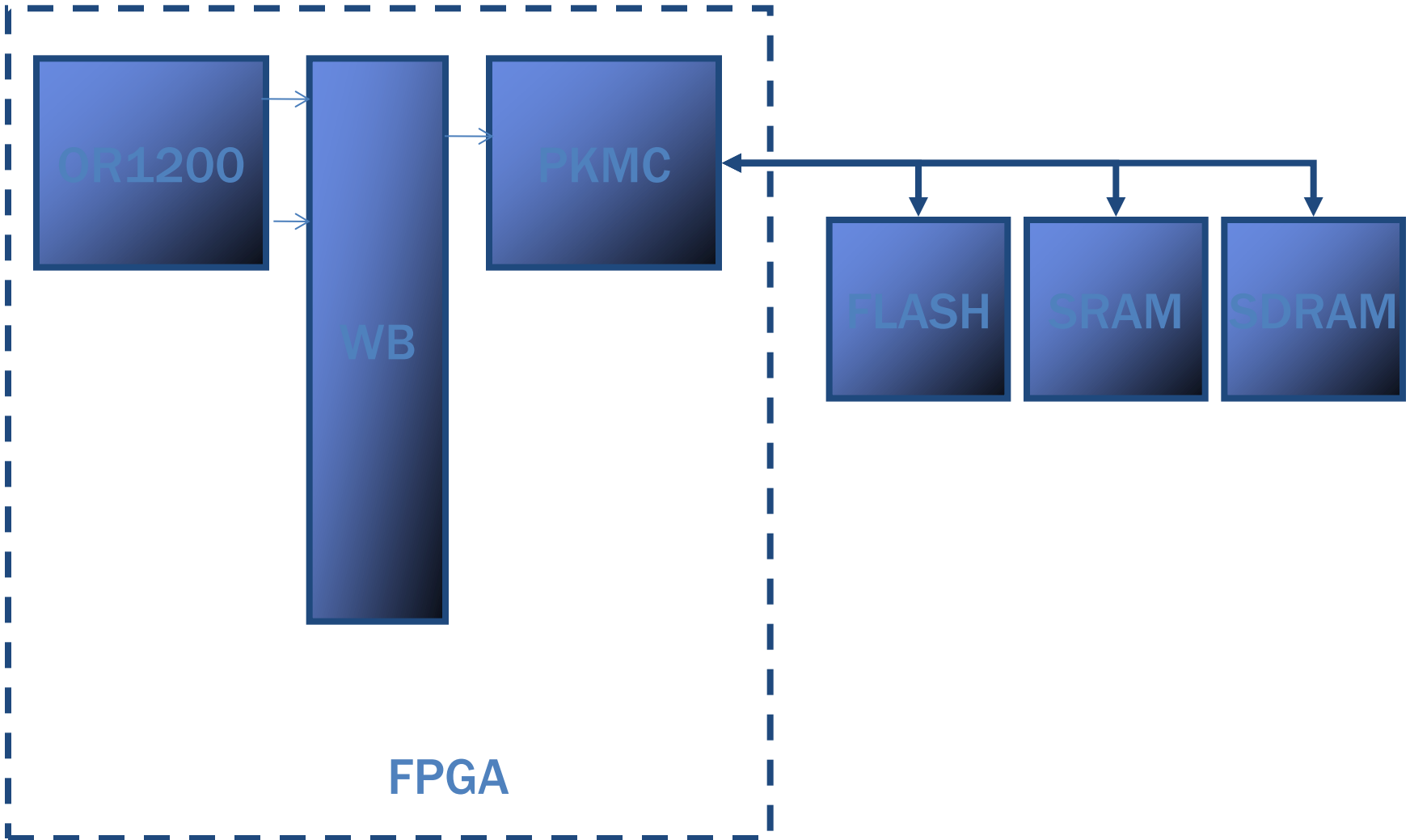
FLASH - Interface

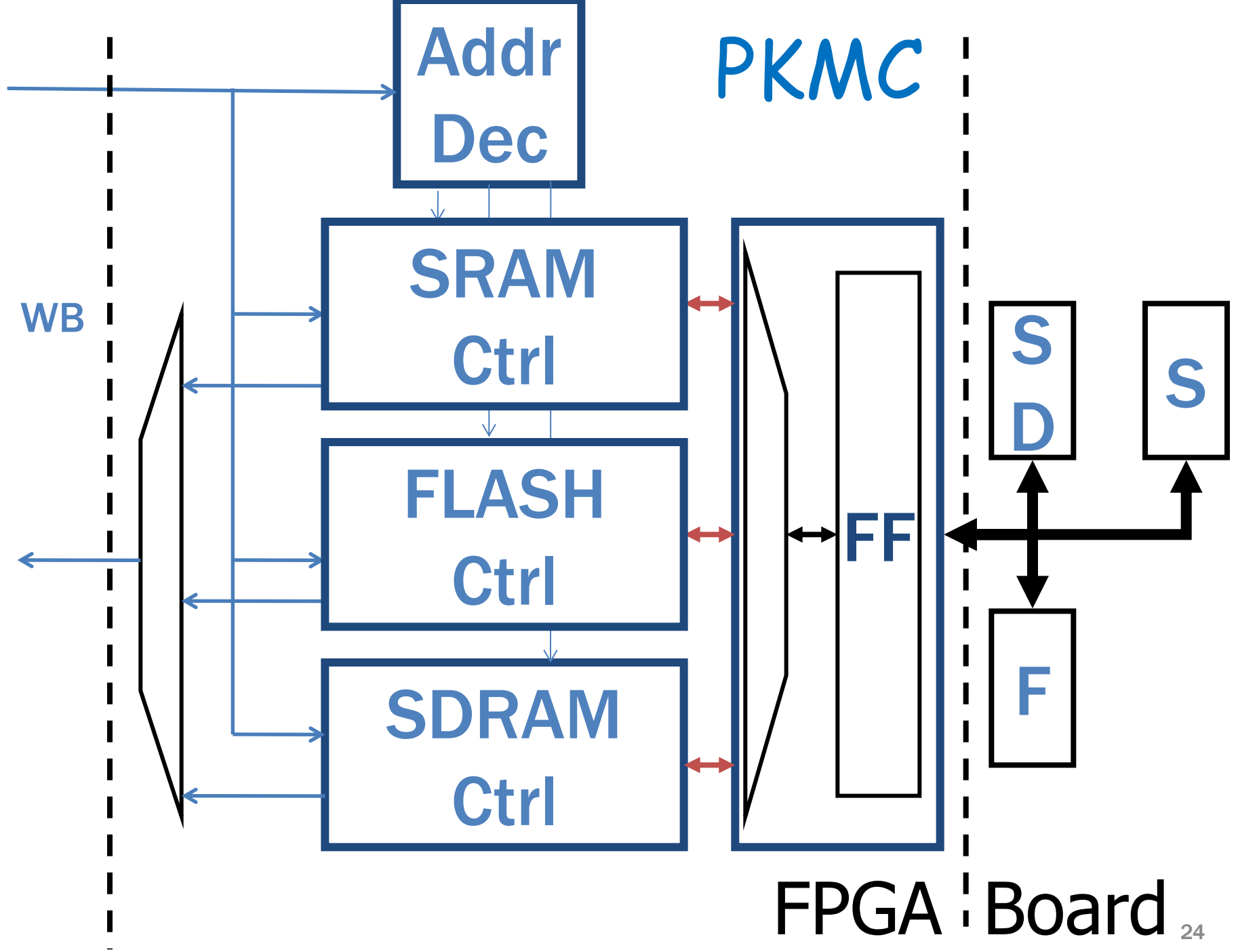
- Looks like SRAM
 - Read
- Write commands
 - Unlock a block
 - Erase a block
 - Program a block
- Contains uCLinux kernel +file system

FLASH - Cell

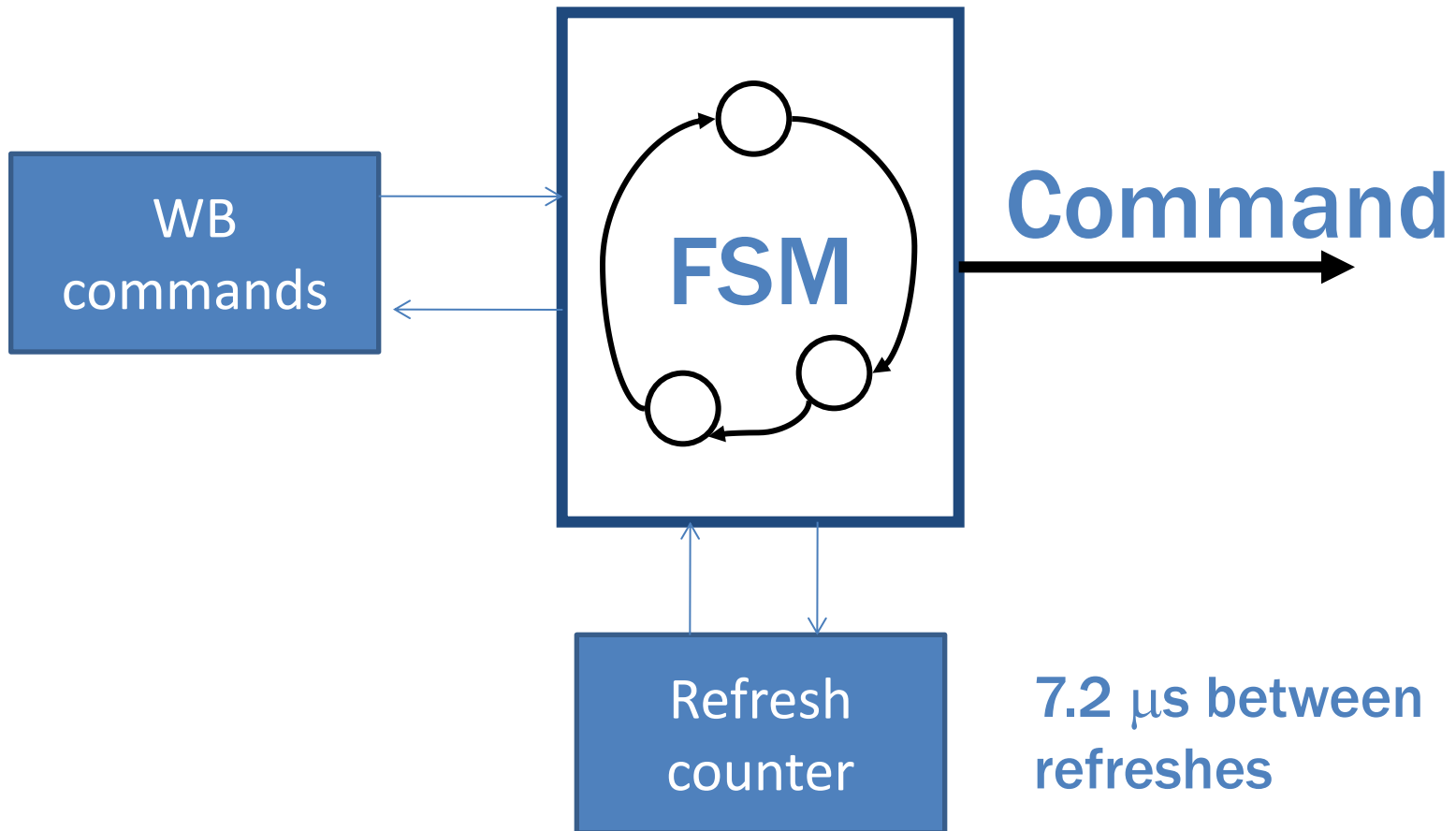


System Overview



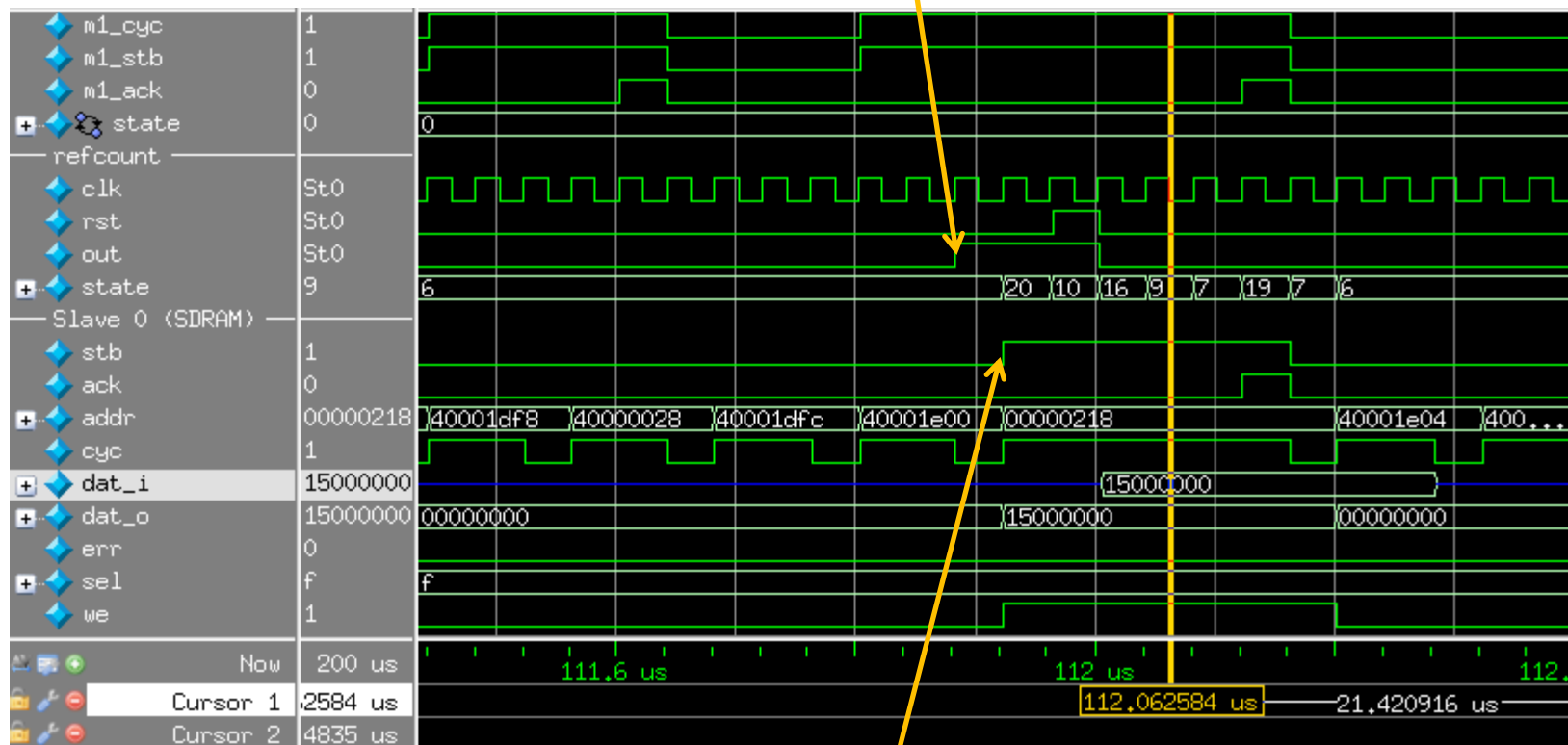


SDRAM Controller Internals



Refresh cycle

Start of refresh cycle



Start of WB cycle

Lab 4

A Custom Instruction

Huffman Encoding/Decoding

1) After Q (for instance):

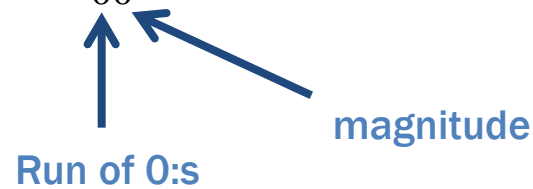
$$\begin{bmatrix} 22 & 12 & 0 & -12 & 0 & 0 & 0 & 0 \\ 0 & 0 & -8 & 0 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

2) After zig-zag:

22
12
0 4
0 0 -12
-8
00000000000000000000
00000000000000000000
000000000000000001
0 ... 0

3) After RLE

value	raw bits
05	10110
04	1100
13	100
24	0011
04	0111
F0	
F0	
D1	1
00	



4) After Huffman

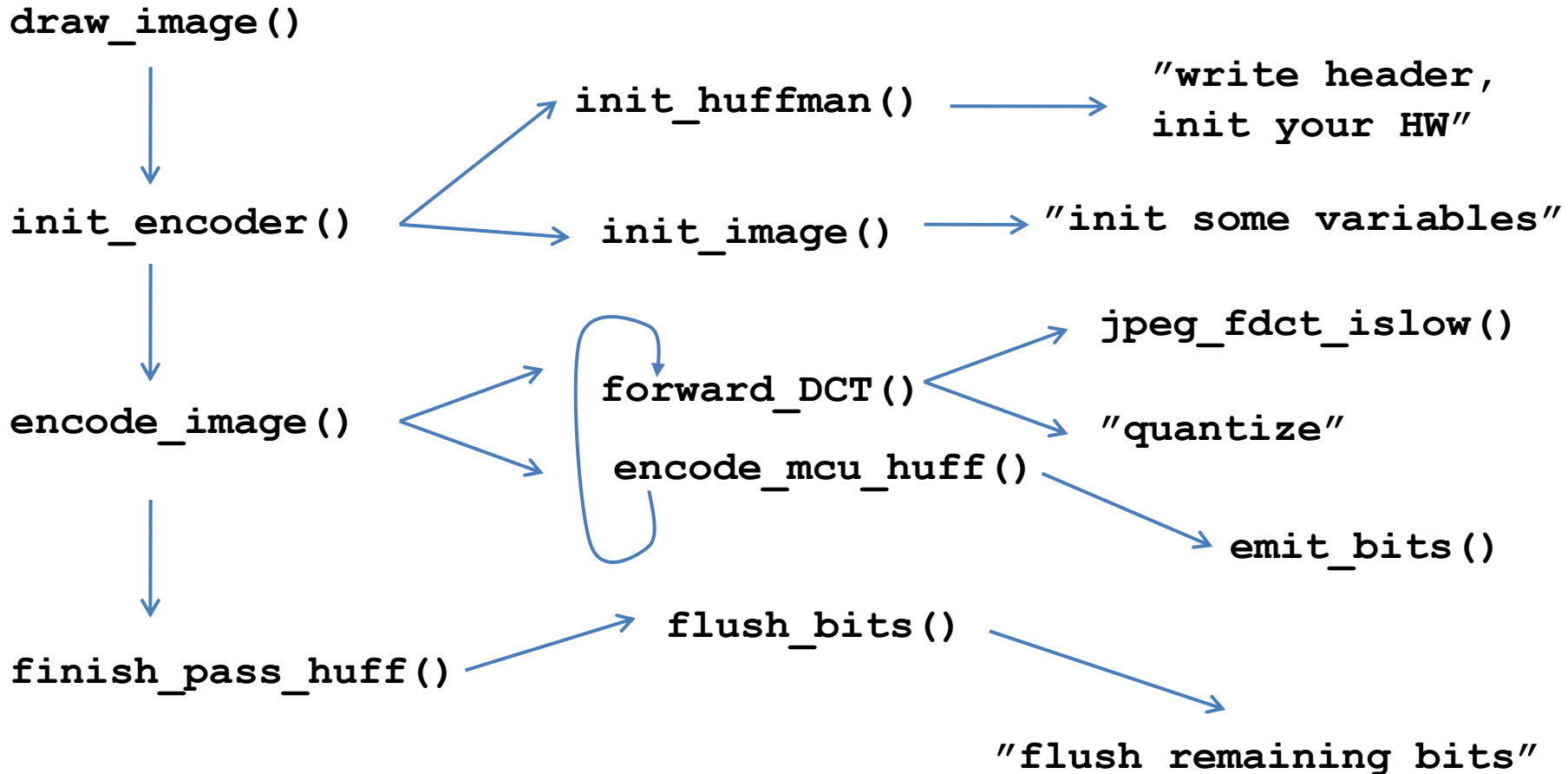
The values (bytes) are encoded by table lookup

Huffman in JFIF

- Output: 1 – 16 bits
- Encodes bytes
- 2 tables used
 - Y DC (differential)
 - Y AC

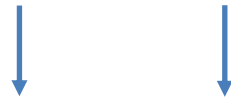
jpegfiles

jpegtest.c, jcdctmgr.c, jdct.c, jchuff.c



Emit_bits()

Max 16



```
static void emit_bits (unsigned int code, int size)
{
    unsigned int startcycle;

    new_put_buffer = (int) code;

    // Add new bits to old bits. If at least 8 bits then write a char to buffer,
    // save the rest until we get more bits.

    new_put_buffer &= (1<<size) - 1;           /* mask off any extra bits in code */
    current_buffer_bit += size;                /* new number of bits in buffer */
    new_put_buffer = new_put_buffer << (24 - current_buffer_bit); /* align incoming bits */
    new_put_buffer = new_put_buffer | old_put_buffer; /* and merge with old buffer contents */

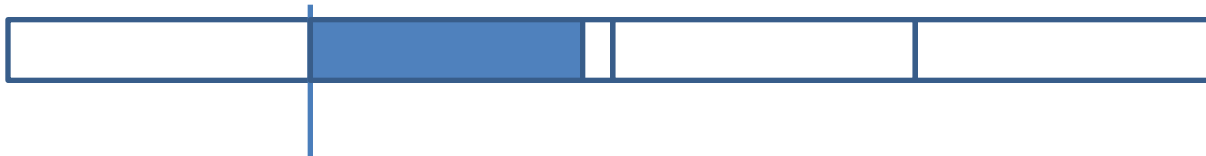
    while (current_buffer_bit >= 8) {
        int c = ((new_put_buffer >> 16) & 0xFF); // Mask out the 8 bits we want
        buffer[next_buffer] = (char) c;
        next_buffer++;
        if (c == 0xFF) { // 0xFF is a reserved code for tags, if we get image data
            buffer[next_buffer] = 0x00; // with an FF value it has to be followed by 0x00.
            next_buffer++;
        }
        new_put_buffer <<= 8;
        current_buffer_bit -= 8;
    }
    old_put_buffer = new_put_buffer; /* update state variables */
}
```

When **buffer** is complete, it is written to file

Emit_bits()

old_put_buffer

current_buffer_bit=7



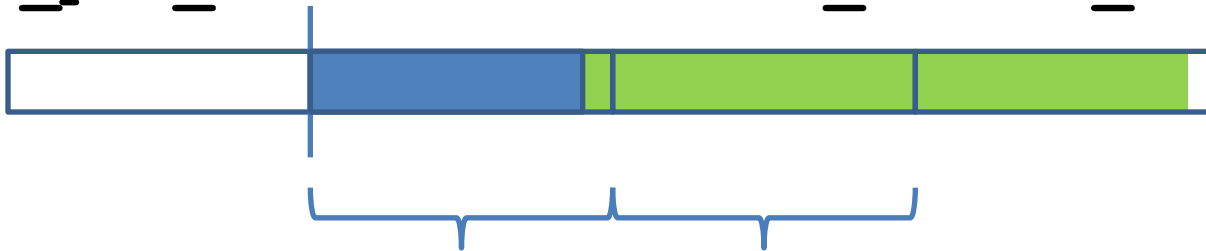
code



size=16

new_put_buffer

current_buffer_bit=23



Write bytes to mem

old_put_buffer

current_buffer_bit=7



Adding an Instruction

1. Instruction Selection
2. Hardware modification
3. (Assembler modification)
4. (Compiler modification)

Instruction Selection

- **l.custx**
 - No operands
- Instructions for 64 bit
 - Not used
 - Assembler can understand
 - **l.sd I(rA), rB**

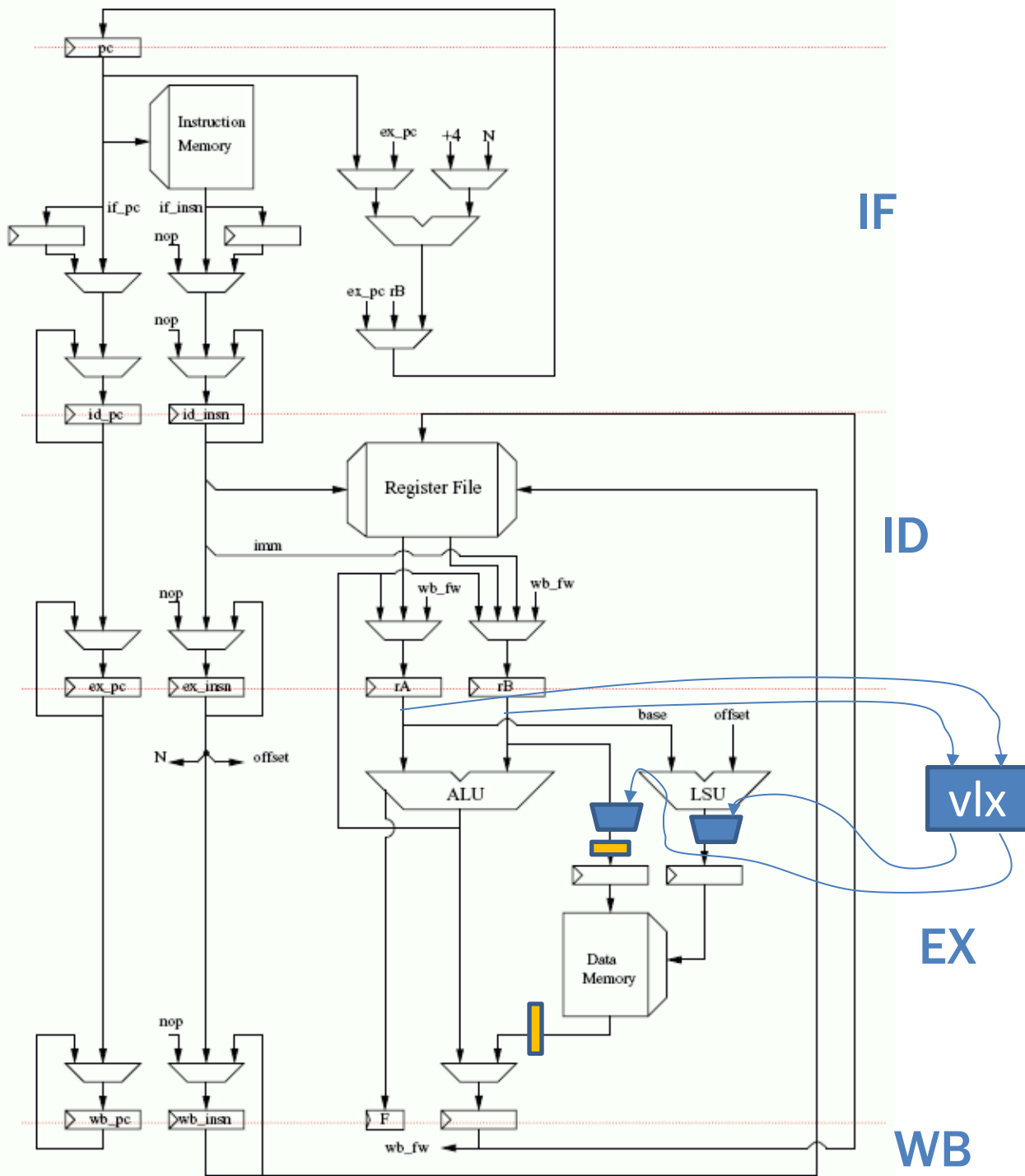


We hijack this instruction

Hardware Modifications

- Instruction decoder modifications
 - Legal instruction
 - or1200_ctrl.v
- Special purpose register
 - New group
 - or1200_sprs.v
- Data path
 - New hardware
 - or1200_lsu.v
 - or1200_vlx_top.v

or1200 Pipeline

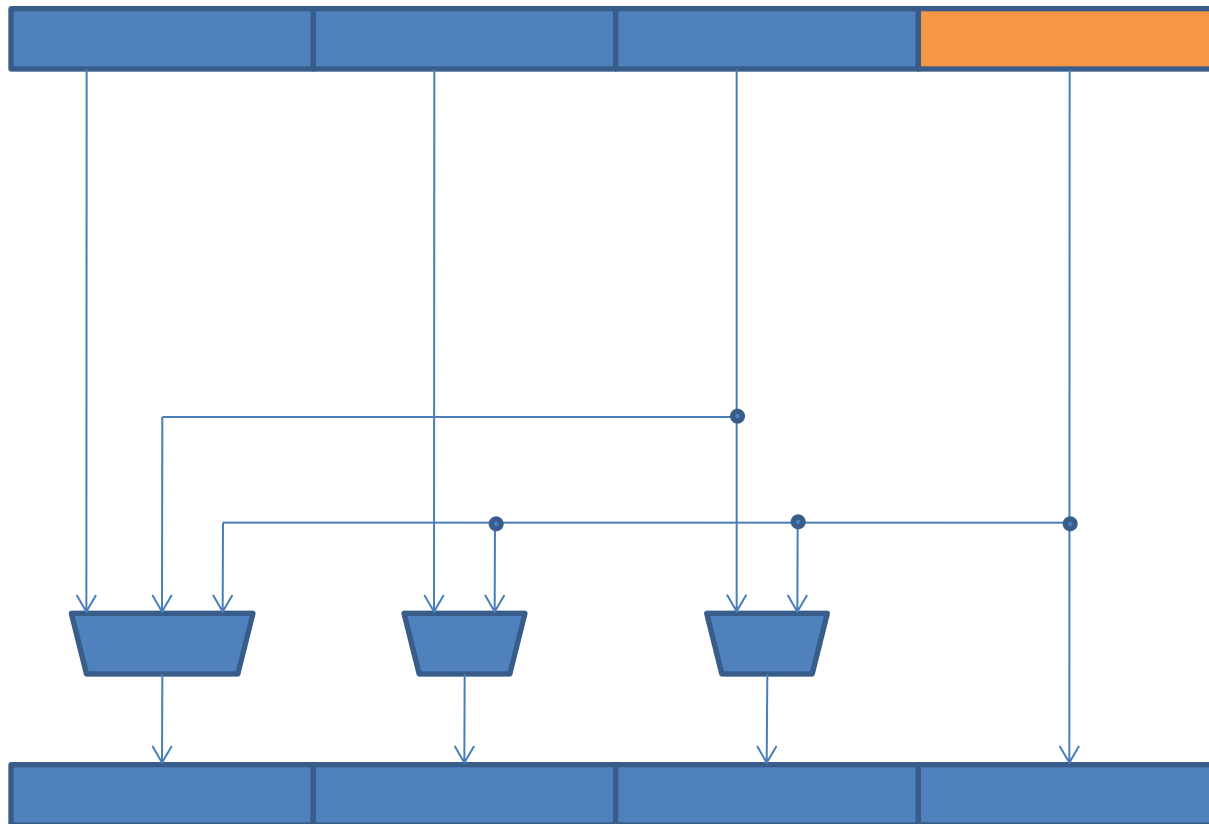


code size
 ↓ ↓
 l.sd (rA), rB

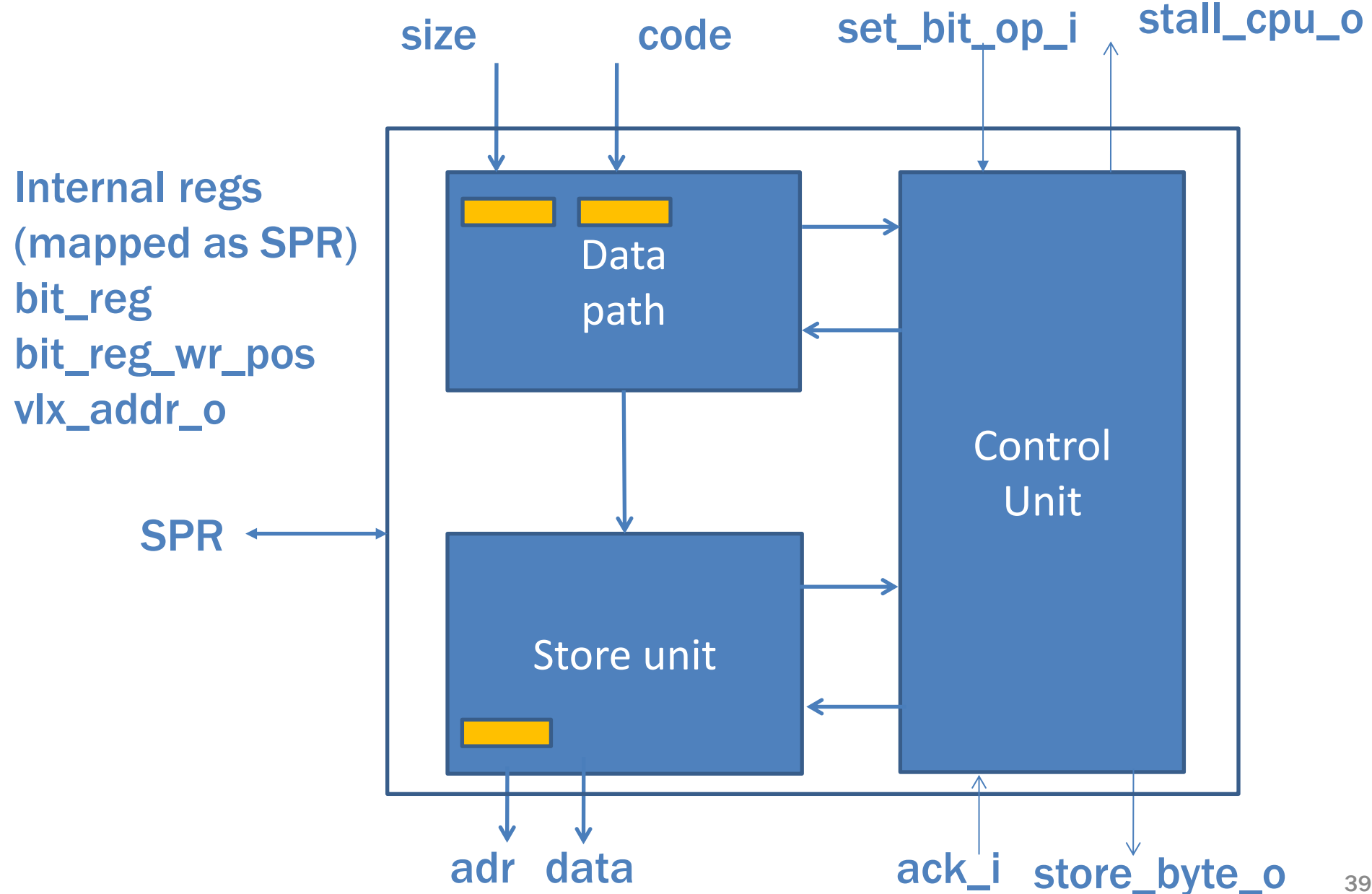
align, reg2mem

align, mem2reg

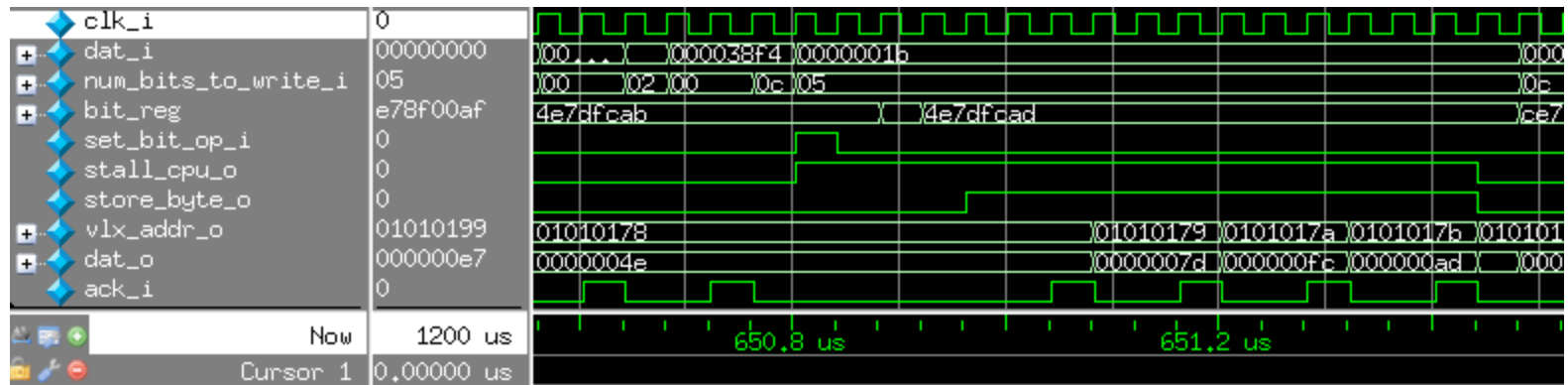
Align reg2mem (for byte write)



Proposed Architecture



An example: PK:s implementation



Inline asm

In jpegfiles insert:

template

```
asm volatile("l.sd 0x0(%0),%1" : : "r"(code), "r"(size));
```

input input

=> code and size will show up at your vx

SPR = special purpose registers

- is an internal bus in the CPU for control and status
- two instructions
 - `mtspr K(rD), rA`
 - `mfspir rD, K(rA)`
- simple bus `spr_cs`, `spr_write`, `spr_addr`, `spr_dat_i`, `spr_dat_o`
- Connect 3 registers to SPR bus!

```
// copy from address register to variable pos
```

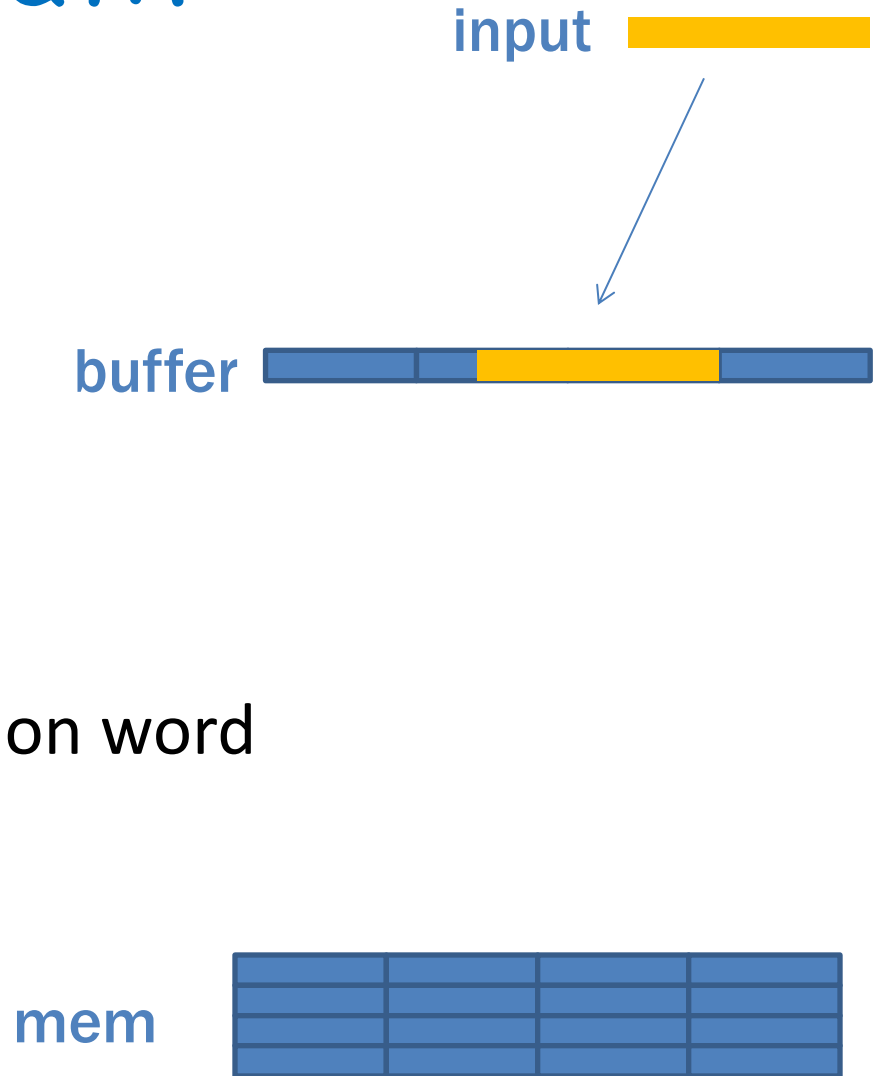
```
asm volatile("l.mfspr %0,%1,0x2" : "=r" (pos) : "r" (0xc000) );
```

```
// zero bit counter
```

```
asm volatile("l.mtspr %0,%1,0x1" : "=r" (0xc000) : "r" (0x0) );
```

Data Path

- Fill buffer
 - One bit / clock cycle
 - All bits at once
- Write to mem
 - One byte
 - One 32 bit word, must be on word boundaries (multiple of 4)



Control

- May not be needed
- May be an FSM

Store Unit

- Stores the data
- 0xFF stored as 0xFF00
 - Jpeg markers
- Only byte alignment!
 - Parallel stores faster

Software

- New Assembler
 - Easy
- New Compiler
 - Hard problem for complex instructions
 - Compiler knows functions
- C
 - Inline Assembler

Instruction Usage

```
unsigned char* sb_get_buff_pos(void)
{
    unsigned char* pos;
    asm volatile("l.mfspr %0,%1,0x2" : "=r" (pos) : "r" (0xc000) );
    return pos;
}
```

output

00000250 <_sb_get_buff_pos>:

250:9c	21	ff	fc	l.addi r1,r1,0xfffffffffc
254:d4	01	10	00	l.sw 0x0(r1),r2
258:9c	41	00	04	l.addi r2,r1,0x4
25c:a9	60	c0	00	l.ori r11,r0,0xc000
260:b5	6b	00	02	l.mfspr r11,r11,0x2
264:84	41	00	00	l.lwz r2,0x0(r1)
268:44	00	48	00	l.jr r9
26c:9c	21	00	04	l.addi r1,r1,0x4