TSTE12 Design of Digital Systems Lecture 7

Kent Palmkvist

LINKÖPING UNIVERSITY

Agenda

- Practical issues
- Algorithm level design
  - Larger models
  - Time multiplexing
- RTL level models
  - Control units
- Gate level models

TSTE12 Deadlines Y,D,ED

- Initial version design sketch and project plan tuesday 12 September
- Weekly meetings should start
  - Internal weekly meeting with transcript sent to supervisor
- Lab 2 soft deadline Wednesday 13 September at 21.00
  - Lab 2 results will be checked after project end

TSTE12 Deadlines MELE, erasmus

- First project meeting no later than today Monday 11 September
- Tuesday 12 September: First version of requirement specification
- Wednesday 13 September 21.00: Lab 1 deadline
  - Pass required to be allowed continued project participation

Handin (homework), Individual!

- 1st handin published today Monday 11 September
  - Deadline Monday 18 September 23:30
- Use only plan text editor (emacs, vi, modelsim or similar) for code entry.
- Solve tasks INDIVIDUALLY
- Submit answers using Lisam assignment function
  - 4 different submissions for code, one for each code task
  - 1 submission for all theory question answers
- Use a special terminal window when working with handins

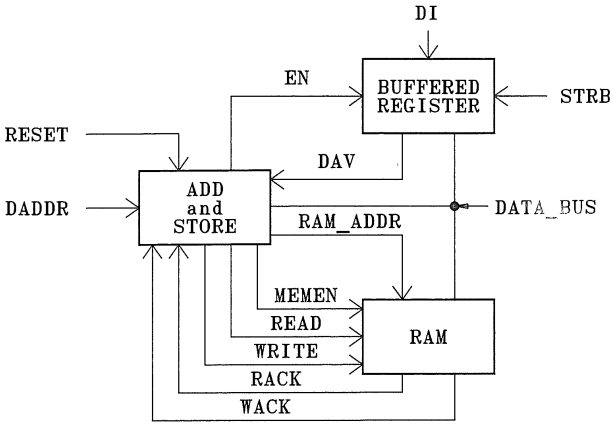
  module load TSTE12 ; TSTE12handin

Algorithm level design

- Focus on functions at high abstraction level
  - Subsystems
  - Algorithms to use
- Ignore timing, datapaths etc.

Bigger system example

- 3 units, Add and Store, Buffered Register, RAM
- Predefined sequence
  - Store value in buffer
  - move buffer value to add and store unit
  - Read memory content and add to input value
  - Write back result to memory



Bigger system example, cont.

- DATA_BUS driven by all modules
  - Requires a resolution function
  - Preset to ZZZ to get a useful value (X will always give X)
  - Note this is done in the entity! Reason: inout => Driver on the entity

```
entity RAM is
  generic(RDEL,DISDEL,ACK_DEL,ACK_PW: TIME);
   port(DATA: inout BUS1(7 downto 0):="ZZZZZZZZ";
        ADDR: in MVL4_VECTOR(4 downto 0);
        RD,WRITE,CS: in MVL4;
        RACK,WACK: out MVL4);
end RAM;
```

Simple RAM model

- Use of MVL4 => use drive and sense functions

```
architecture SIMPLE of RAM   is
begin
 MEM: process (CS,RD,WRITE)
  type MEMORY is array(0 to 31) of MVL4_VECTOR(7 downto 0);
  variable MEM: MEMORY:= (others => (others  => '0'));
 begin
   if CS = '1' then
     if RD = '1' then
       DATA <= DRIVE(MEM(INTVAL(ADDR))) after RDEL;
       RACK <= '1' after ACK_DEL,
               '0' after ACK_DEL + ACK_PW;
     elsif WRITE = '1' then
       MEM (INTVAL(ADDR)):= SENSE(DATA,'1');
       WACK <= '1' after ACK_DEL, '0' after ACK_DEL+ACK_PW;
     end if;
   else
     DATA <= "ZZZZZZZZ"  after DISDEL;
   end if;
 end process MEM;
end SIMPLE;
```


Bigger example, cont.

- The RAM-model uses an aggregate to initialize all elements to zero
- ADD and Store is a form of a state machine
  - Go through a sequence step by step
  - Execute some function in each step
  - Each step ends in a wait
- Divide system into datapath and control
- Clock generation as earlier (loop with run)

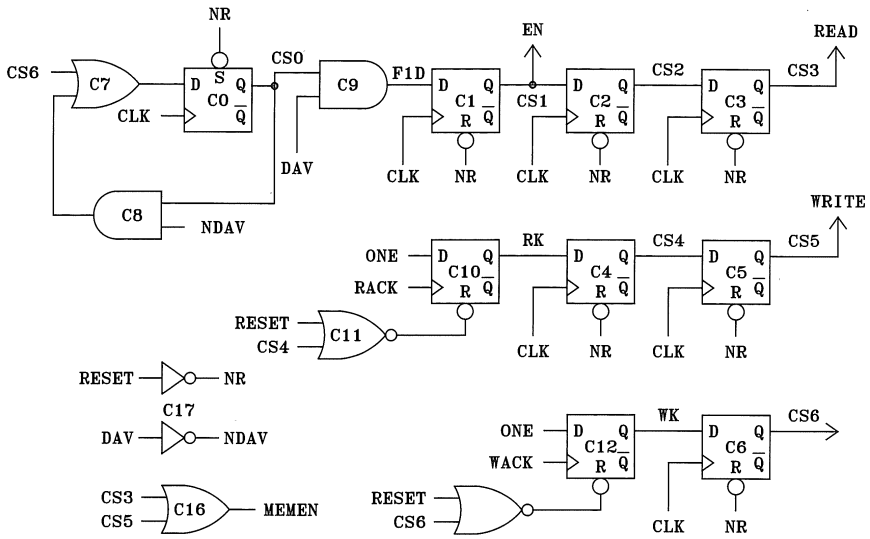
Bigger example, cont.

```
CON: process
  variable DATA_REG:
             MVL4_VECTOR(7 downto 0);
 begin
   if RESET = '1' then           --CS0
     DATA <= "ZZZZZZZZ"after DIS_DEL;
   end if;
   wait on DAV until DAV = '1';
----------------------------------------
   EN <= '1' after CON_DEL;      --CS1
   wait for CLK_PER;
----------------------------------------
   EN <= '0' after CON_DEL;
   DATA_REG := SENSE(DATA,'1');  --CS2
   wait for CLK_PER;
----------------------------------------
   MADDR <=  DADDR after MA_DEL;
   MEMEN <= '1' after CON_DEL;   --CS3
   READ <= '1' after CON_DEL;
   wait on RACK until RACK ='1';
----------------------------------------
   DATA_REG :=
      ADD8(SENSE(DATA,'1'),DATA_REG);
   READ <= '0'after CON_DEL;
   MEMEN <= '0'after CON_DEL;    --CS4
   wait for CLK_PER;
----------------------------------------
   DATA <= DRIVE(DATA_REG) after DO_DEL;
   WRITE <= '1'after CON_DEL;
   MEMEN <= '1'after CON_DEL;    --CS5
   wait on WACK until WACK ='1';
----------------------------------------
   WRITE <= '0'after CON_DEL;
   MEMEN <= '0'after CON_DEL;    --CS6
   DATA <= "ZZZZZZZZ" after DIS_DEL;
   wait for CLK_PER;
end process CON;
```


Control state machine

- Hardware aspects on the control machine
  - Wait can not be used in synthesis
  - Use a manual direct translation technique
- One-hot encoding
  - Simple and straight forward
  - Suitable for FPGA implementation
  - Low complexity decoding of state

One-hot encoded controller



Control step classes

- Automatic increase
  - From step c_i to c_{i+1} after some time
- Handshake
  - Wait for DAV, CS1 => EN = 1, Buffer resets DAV when EN = 1
- Asynchronous stepping
  - CS3 to CS4: Wait for external RACK edge, RACK may be shorter than 1 clock period!

Hardware vs Behavioral model

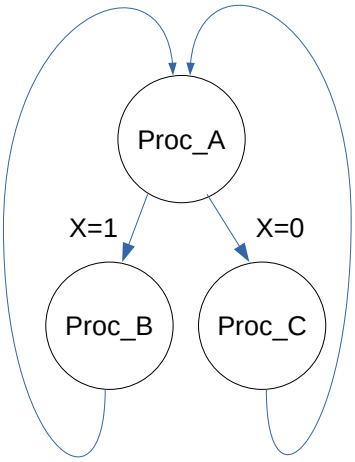
- Important to have same behavior of hardware and VHDL model
- Reset behavior is different
  - The model only checks for reset in CS0
  - Hardware checked reset everywhere
  - Different behavior between model and HW! Bad.
  - Add reset check in every control step

Why single clock domainn

- Reset problems
  - Even single clock domain should synchronize asynchronous reset inputs
  - Must guarantee that whole circuit releases from reset at the same time
- Communication problems
  - Possible race between data and clock
  - Metastability

Control branching

- Want a control machine able to handle conditional execution of sequences
  - Similar to hierarchical finite state machine (FSM)
  - VHDL cannot do jumps, only breaking loops
- Working implementation approach
  - Sequences in individual processes
  - Check at end of process which process to start next
- Output signals from state machine
  - Require Resolution function as assignment done in multiple processes (need to turn off non-active processes)



Each processes should control signal S

Control branching using multiple processes

```
architecture TWO of WAIT_STEPS is

 signal TRIGGERB,TRIGGERBA,
        TRIGGERC,TRIGGERCA: DOT1 := '0';
 signal SINT: RINTEGER register;

begin

A: process
  begin
    SINT <= null;
   wait on RUN,TRIGGERBA,TRIGGERCA until RUN = '1';
    SINT <= 0;                    ---Step 0
   wait for CLK_DEL;
    SINT <= 1;                    ---Step 1
   wait for CLK_DEL;
    SINT <= null;
    if X = '1' then
      TRIGGERB <= not(TRIGGERB);
    else
      TRIGGERC <= not(TRIGGERC);
    end if;
  end process A;

B: process
  begin
    SINT <= null;
   wait on TRIGGERB;
    SINT <= 2;                    ---Step 2
   wait for CLK_DEL;
    SINT <= 3;                    ---Step 3
   wait for CLK_DEL;
    SINT <= null;
    TRIGGERBA <= not(TRIGGERBA);
  end process B;

C: process
  begin
    SINT <= null;
   wait on TRIGGERC;
    SINT <= 4;                    ---Step 4
   wait for CLK_DEL;
    SINT <= 5;                    ---Step 5
   wait for CLK_DEL;
    SINT <= null;
    TRIGGERCA <= not(TRIGGERCA);
  end process C;
  S <= SINT;

end TWO;
```

Time multiplexing

- Problem: Multiple processes driving one signal
  - Multiple drivers (one for each process)
  - Want to enable separate drivers at non-overlapping intervals
  - Assigned signal value should keep the value even after driver enable removed (memory function)
  - Use signal type containing a resolution function
- Remember: This is NOT for synthesis

Time Multiplexing, non-working

- Two-phase clock
  - One pulse each alternating
- Resolved output signal Z
  - Allows multiple assignment
- Problem
  - Z = ’X’ when PH_TWO assign ’1’
  - Assignment from PH_ONE will not turn off
  - Each driver always outputs last assigned value

```
entity TIME_MUX is
  generic(DEL1,DEL2: TIME);
  port(PHASE_ONE,PHASE_TWO: in MVL4;
       Z: out DOTX := '0');
end TIME_MUX;

architecture PROCESS_IF_0 of TIME_MUX is
begin

  PH_ONE:process(PHASE_ONE)
  begin
    if PHASE_ONE = ’1’ then
      Z <= '0' after DEL1;
    end if;
  end process;

  PH_TWO:process(PHASE_TWO)='1')
  begin
    if PHASE_TWO = ’1’ then
      Z <= '1' after DEL2;
    end if;
  end PH_TWO;

end PROCESS_IF_0;
```

Time multiplexing, cont.

- Ordinary processes
  - Assignment of null disables driver
  - Keyword register in signal declaration
    - defines what happen when all driver are null

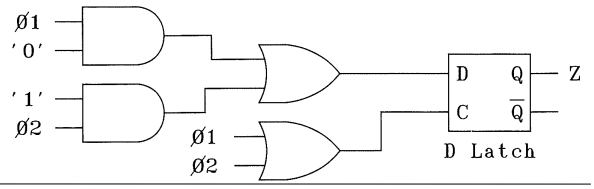
```
architecture PROC_NULL of TIME_MUX is
  signal ZINT: DOTX register;
begin

 PH_ONE: process(PHASE_ONE)
 begin
  if PHASE_ONE = '1' then
    ZINT <= '0' after DEL1;
  else
    ZINT <= null;
  end if;
 end process PH_ONE;

 PH_TWO: process(PHASE_TWO)
 begin
  if PHASE_TWO = '1' then
    ZINT <= '1' after DEL2;
  else
    ZINT <= null;
  end if;
 end process PH_TWO;

 Z <= ZINT;

end PROC  NULL:
```



Time multiplexing, cont.

- Use bus instead of register
  - Default resolution function used if not driven
    - Resolution function recieves an empty input
  - MVL4 resolution function start with ’Z’ value!

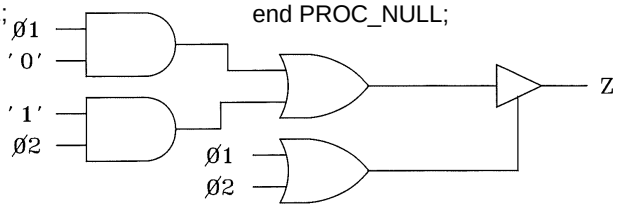
```
architecture PROC_NULL of TIME_MUX is
  signal ZINT: DOTX bus;
begin

 PH_ONE: process(PHASE_ONE)
 begin
  if PHASE_ONE = '1' then
    ZINT <= '0' after DEL1;
  else
    ZINT <= null;
  end if;
 end process PH_ONE;

 PH_TWO: process(PHASE_TWO)
 begin
  if PHASE_TWO = '1' then
    ZINT <= '1' after DEL2;
  else
    ZINT <= null;
  end if;
 end process PH_TWO;

 Z <= ZINT;

end PROC_NULL;
```



Time multiplexing, cont.

- Without the use of Register/Bus.
- Separate signals, check 'QUIET to find active assignment

```
entity TIME_MUX is
  generic(DEL1,DEL2: TIME);
  port(PHASE_ONE,PHASE_TWO: in
MVL4;
       Z: buffer MVL4);
end TIME_MUX;

architecture QUIET_MUX of TIME_MUX is
  signal  PH1,PH2,Z1,Z2: MVL4;
begin

  PH_ONE: process(PHASE_ONE)
  begin
    if PHASE_ONE = '1' then
      Z1 <= '0' after DEL1;
    end if;
  end process PH_ONE;

  PH_TWO: process(PHASE_TWO)
  begin
    if PHASE_TWO = '1' then
      Z2 <= '1' after DEL2;
    end if;
  end process PH_TWO;

  Z <= Z1 when not Z1'quiet else
       Z2 when not Z2'quiet else
       Z;

end QUIET_MUX;
```


Register transfer level (RTL)

- At this level can the following aspects be analysed
  - Compare timing between different units at register level
    - Delay in subfunctions, etc.
  - Resource allocation
    - Number of buses, registers, processing elements etc.
  - Scheduling (when to perform an operation)
  - Control structure (e.g., microcoded control units)
  - Bus design

Difference between behavior and dataflow descriptions

- Behavioral model
  - System with two registers and an adder
  - Behavior description does not indicate how operations are performed
  - Command selects operation
  - Only signals that corresponds to saved data

```
entity REG_SYS is
  port(C: in BIT;
       COM: in BIT_VECTOR(0 to 1);
       INP: in BIT_VECTOR(0 to 7));
end REG_SYS;

architecture ALG of REG_SYS is
  signal R1,R2: BIT_VECTOR(0 to 7);
begin

  process(C)
  begin
    if C='1' then
      case COM is
        when "00" => R1 <= INP;
        when "01" => R2 <= INP;
        when "10" => R1 <= ADD8(R1,R2);
        when "11" =>
           R1 <= ADD8(R1,INC8(not(R2)));
      end case;
    end if;
  end process;

end ALG;
```


Behavioral vs Dataflow, cont.

- Dataflow model
  - More signals (many for communication)
  - Operations are registers, multiplexes, or arithmetic/logic operations
  - Global decoding using signals D00 to D11
  - Corresponds to a data flow graph

Behavioral vs Dataflow, cont.

```
architecture DF1 of REG_SYS is
  signal MUX_R1,R1,R2,R2C,R2TC,MUX_ADD,SUM:
         BIT_VECTOR(0 to 7);
  signal D00,D01,D10,D11,R1E: BIT;

begin
  D00 <= not COM(0) and not COM(1);
  D01 <= not COM(0) and COM(1);    ---Command Decoder
  D10 <= COM(0) and not COM(1);
  D11 <= COM(0) and COM(1);
  MUX_R1  <= SUM when D00 = '0' else INP; --Reg 1 Mux
  R1E <= D00 or D10 or D11;

R1_REG: process(C)    -- Register 1
  begin
    if (C='1') and C'EVENT) then
      if (R1E = '1') then
        R1 <= MUX_R1;
      end if;
    end if;
  end process;
```

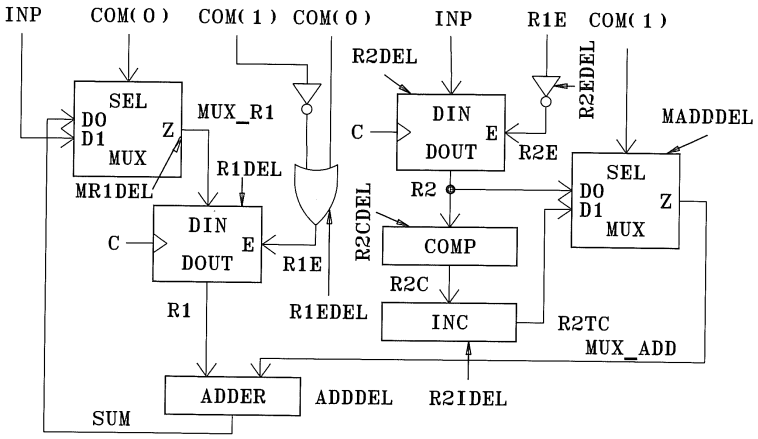
```
R2_REG: process(C)       --Register 2
  begin
    if (C='1') and C'EVENT) then
      if (D01 = '1') then
        R2 <= INP;
      end if;
    end if;
  end process;

  R2C <= not R2;                          ---Complement
  R2TC <= INC8(R2C);                      ---Increment
  MUX_ADD <= R2TC when D11 = '1' else R2; ---Adder Mux
  SUM <= ADD8(R1,MUX_ADD);                ---Adder

end DF1;
```


Described Dataflow implementation

- One-to-one mapping



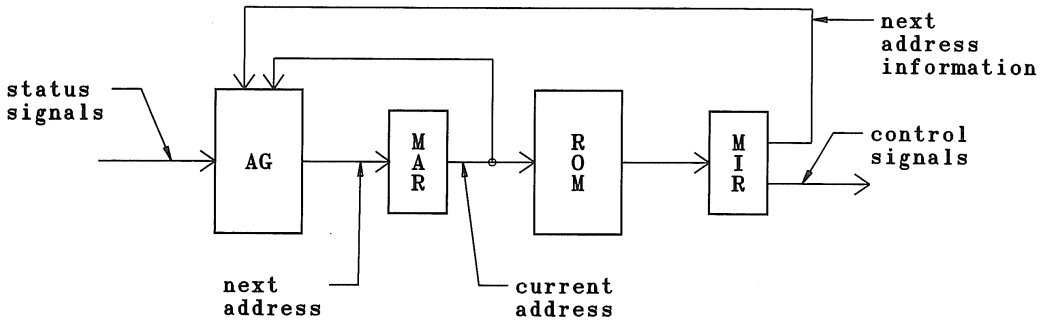

```
  MUX_R1  <= SUM when D00 = '0'
                 else INP;
  R1E <= D00 or D10 or D11;
R1_REG: process(C) begin
  if (C='1') and C'EVENT and
      (R1E='1') then
    R1 <= MUX_R1;
  end if; end process;
R2_REG: process(C) begin
  if (C='1') and C'EVENT and
     (D01='1') then
    R2 <= INP;
  end if; end process;
  R2C <= not R2;
  R2TC <= INC8(R2C);
  MUX_ADD <= R2TC when D11 = '1'
             else R2;
  SUM <= ADD8(R1,MUX_ADD);
```

Control units

- Hard wired
  - Moore (output only dependent on state)
  - Mealy (output dependent on state and input)
  - Fast
  - Custom designed
- Microcoded
  - Cheap
  - Standardized (easy to reuse)

Microcoded control unit

- General structure
  - AG = Adress generator
  - MAR = Memory Adress Register
  - MIR = Memory Instruction register

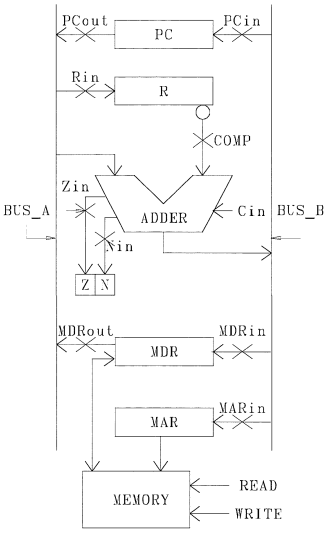


Microcoded control unit

- Advantages
  - Easy to create a generic design
  - Only ROM contents needs to be replaced
  - Easy to change existing design
  - Short design time (low design cost)
  - May use compiler to create ROM contents
- Drawbacks
  - Slower in many cases (ROM must be read)
    - Only Moore type of controllers
  - Small controllers are more expensive due to extra register and ROM
  - Must be designed for worst case regarding required features

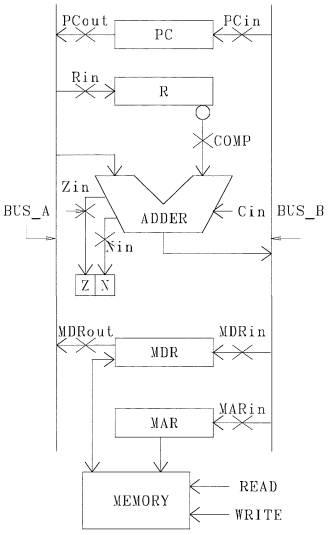
Microcoded control unit, example

- Controller for an extremly small RISC processor
  - 4 register (PC, R, MDR, MIR)
  - 1 subtraction unit
  - Some multiplexers and busses
  - Use the same add unit both for instruction operation and PC update
  - Cost: 9 clock cycles per instruction
- Only one instruction: subtract with branch on negative result
  - 3 byte instruction
    - 1st operand address, 2nd operand address, branch address



Controller for structure

- One instruction execution
  0. PCout,Zin,MARin,READ,ZEND
  1. MDRout,MARin,READ
  2. MDRout,Rin
  3. Pcout,Cin,PCin,MARin,READ
  4. MDRout,MARin,READ
  5. MDRout,COMP,Cin,Nin,MDRin,WRITE
  6. PCout,Cin,PCin,MARin,READ
  7. PCout,Cin,PCin,NNEND
  8. MDRout,Pcin
- 2 loops, 0-7 or 0-8



Control with two jumps, microcoded

- All control steps described in a ROM table
- Easy to understand
- Easy to redesign

```
-------- The microinstructions ROM  ----------------------
ROM: process(C)
  type SQ_ARRAY is array(0 to 8,0 to 8) of BIT;
  constant MEM : SQ_ARRAY:=
-- 0       1    2    3     4     5    6    7     8   COLUMN
MDR_OUT,MAR_IN,N_IN,R_IN,PC_IN,ZEND,C_IN,WRITE,NNEND,micins
(('0',    '1', '0', '0',  '0',  '1', '0', '0',  '0'), --0
 ('1',    '1', '0', '0',  '0',  '0', '0', '0',  '0'), --1
 ('1',    '0', '0', '1',  '0',  '0', '0', '0',  '0'), --2
 ('0',    '1', '0', '0',  '1',  '0', '1', '0',  '0'), --3
 ('1',    '1', '0', '0',  '0',  '0', '0', '0',  '0'), --4
 ('1',    '0', '1', '0',  '0',  '0', '1', '1',  '0'), --5
 ('0',    '1', '0', '0',  '1',  '0', '1', '0',  '0'), --6
 ('0',    '0', '0', '0',  '1',  '0', '1', '0',  '1'), --7
 ('1',    '0', '0', '0',  '1',  '0', '0', '0',  '0'));--8
begin
  MDR_OUT  <=   MEM(INTVAL(C),0) after ROM_DEL;
  MAR_IN   <=   MEM(INTVAL(C),1) after ROM_DEL;
  N_IN     <=   MEM(INTVAL(C),2) after ROM_DEL;
  R_IN     <=   MEM(INTVAL(C),3) after ROM_DEL;
  PC_IN    <=   MEM(INTVAL(C),4) after ROM_DEL;
  ZEND     <=   MEM(INTVAL(C),5) after ROM_DEL;
  C_IN     <=   MEM(INTVAL(C),6) after ROM_DEL;
  WRITE    <=   MEM(INTVAL(C),7) after ROM_DEL;
  NNEND    <=   MEM(INTVAL(C),8) after ROM_DEL;
end process ROM;
```

URISC controller, Mealy

- Inclear sequence
- Hard to modify
- Faster

```
--Hard Wired Control Unit
--Decoder
--First Stage Decoding
ST0 <= not C(2) and not C(1) and not C(0) after AND_DEL;
ST1 <= not C(2) and not C(1) and C(0) after AND_DEL;
ST2 <= not C(2) and C(1) and not C(0) after AND_DEL;
ST3 <= not C(2) and C(1) and  C(0) after AND_DEL;
ST4 <= C(2) and not C(1) and not C(0) after AND_DEL;
ST5 <= C(2) and not C(1) and C(0) after AND_DEL;
ST6 <= C(2) and C(1) and not C(0) after AND_DEL;
ST7 <= C(2) and C(1) and C(0) after AND_DEL;
--Second Stage Decoding
ST07 <= ST0 or ST7 after OR_DEL;
ST25 <= ST2 or ST5 after OR_DEL;
ST36 <= ST3 or ST6 after OR_DEL;
ST57 <= ST5 or ST7 after OR_DEL;
ST78 <= ST7 or C(3) after OR_DEL;
--Control Signals
PC_OUT <= (ST07 or ST36) and not C(3)
          after (OR_DEL + AND_DEL);
C_IN <= ST36 or ST57 after OR_DEL;
PC_IN <= ST36 or ST78 after OR_DEL;
MAR_IN <= not(ST25 or ST78) after (OR_DEL + INV_DEL);
MDR_OUT <=not PC_OUT after INV_DEL;
READ <= MAR_IN; COMP <= ST5; N_IN <= ST5; MDR_IN <= ST5;
WRITE <= ST5; R_IN <= ST2; ZIN <= ST0; ZEND <=ST0;
NNEND <= ST7;
```


More on microcoded controllers

- Lecture 11 will cover more details on microcoded controller structures
  - Introduces also lab 3
- Lab 3 includes an example of a microcoded controller structure
  - Controller used to control a user interface and a datapath
  - Y and D program students have seen this approach in computer technology courses
    - Used there for creating machine instruction implementations

Gate level simulation

- All designs will eventually reach the gate level
- Need accuracy to allow check of timing requirements
  - Setup time on flip-flops
  - Clock signals
  - Races, hazards
  - Glitch example (inverter + and with rising edge input)
- Models must be efficient
  - Large number of gates
  - Slow simulation due to accuracy
    - Still much faster than spice simulation

How accurate can a gate model be?

- Example: 2 input OR-gate

```
Entity OR2 IS
  Port (I1, I2 : in bit; O : out bit);
END OR2;
Architecture DELTA_DEL of OR2 IS
BEGIN
 O <= I1 OR I2;
END DELTA_DEL;
Architecture FIXED_DEL OF OR2 IS
BEGIN
  O <= I1 OR I2 after 3 ns;
END FIXED_DEL;
```

```
ENTITY OR2G IS
  Generic (DEL: TIME)M
  Port (I1, I2 : in bit; O : out
    bit);
END OR2G;
Architecture GNR_DEL of
    OR2G IS
BEGIN
  O <= I1 OR I2 after DEL;
END GNR_DEL;
```

Model accuracy

- Models are better and better, but not good enough
  - Multiple timing models required
  - typical delay, max, min
- Want single model, only changing one constant
  - Timing_CONTROL
  - Set one constant to define type of timing (min, max, typical)

Code example

```
use work.TIMING_CONTROL.all;
entity OR2_TV is
  generic(TMIN,TMAX,TTYP: TIME);
  port(I1,I2: in BIT; O: out BIT);
end OR2_TV;

architecture VAR_T of OR2_TV is
begin
  O <= I1 or I2 after T_CHOICE(TIMING_SEL,
                TMIN,TMAX,TTYP);
end VAR_T;
```

```
package TIMING_CONTROL is
  type TIMING is (MIN,MAX,TYP,DELTA);
  constant TIMING_SEL: TIMING := TYP;
  function T_CHOICE(TIMING_SEL: TIMING;
                    TMIN,TMAX,TTYP: TIME)
    return TIME;
end TIMING_CONTROL;

package body TIMING_CONTROL is
  function T_CHOICE(TIMING_SEL: TIMING;
                    TMIN,TMAX,TTYP: TIME)
    return TIME is
  begin
    case TIMING_SEL is
      when DELTA => return 0 ns;
      when TYP => return TTYP;
      when MAX => return TMAX;
      when MIN => return TMIN;
    end case;
  end T_CHOICE;
end TIMING_CONTROL;
```

Additional timing details

- Timing is asymmetric
  - Different rise and fall times
  - Needs modeling

```
entity OR2GV is
  generic(TPLH,TPHL: TIME);
  port(I1,I2: in BIT; O: out BIT);
end OR2GV;

architecture VAR_DEL of OR2GV is
begin
  process(I1,I2)
    variable OR_NEW,OR_OLD:BIT;
  begin
    OR_NEW := I1 or I2;
    if OR_NEW = '1' and OR_OLD = '0' then
      O <= OR_NEW after TPLH;
    elsif OR_NEW = '0' and OR_OLD = '1' then
      O <= OR_NEW after TPHL;
    end if;
    OR_OLD := OR_NEW;
  end process;
end VAR_DEL;
```


Load dependency

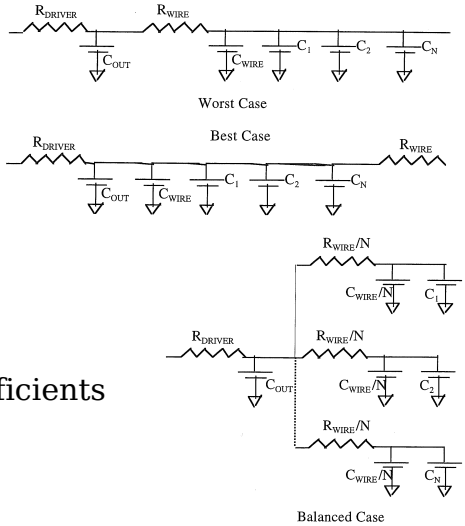
- Every attached gate input slows the output speed
  - Large fan-out
  - Load is gate dependent
    - Number of transistor gates connected
    - Size of transistors on input gate
- Each connection corresponds to a small delay
  - Model each individual input wire delay
  - Gate delay included in output wire delay
- Not good enough still
  - Delay depends on edge slope, temperature, etc.

Common model used in synopsis library compiler

- $D_{TOTAL} = D_I + D_S + D_T + D_C$
- D_I = Intrinsic delay inherent in gate and independent of where/how it is used
- D_S = Slope delay caused by ramp time of the input signal
- D_T Transition delay caused by loading of the output pin (approx R_{driver} (C_{wire}+C_{pin}))
- D_C Connect media delay to an input pin (wire delay).

Different max and min

- Wire delay (D_C) more complicated
  - Worst case
  - Best case
  - Balanced
- Technology library
  - Large amount of information
  - Usually described as tables
  - Sometimes described as polynomial coefficients



Back annotation

- The process of abstraction
  - adding more details to a high level model by analyzing a lower abstraction level model
  - Example: Layout information used to generate timing information in a gate netlist
- Standardized way: SDF
  - Add timing info from layout to gate level
  - Useful for general timing requirements and properties)
  - Delays module path, device, interconnect, and port

SDF file format

- Timing checks: setup, hold, recovery, removal, skew, width, period, and no change
- Timing constraints: path, skew, period, sum, and diff
- Each trippel defines min, typical, and max delay
  - One for positive edge
  - One for negative edge

```
(CELL
  (CELLTYPE "DFF")
  (INSTANCE top/b/c)
  (DELAY
    (ABSOLUTE
      (IOPATH (posedge clk) q (2:3:4) (5:6:7))
      (PORT clr (2:3:4) (5:6:7))
    )
  )
  (TIMINGCHECK
    (SETUPHOLD d (posedge clk) (3:4:5) (-1:-1:-1))
    (WIDTH clk (4.4:7.5:11.3))
  )
)
```

SDF File format, cont.

- Design/instance-specific or type/library-specific data
- Timing environment:
- intended operating timing environment
- Scaling, environmental, and technology parameters
- Incremental delay builds on the previous models timing by adding/subtracting timing information
- Absolute replaces timing information

Gate models of increasing complexity

- Creating accurate library models is time consuming
- Delay, timechecks etc. can be done in many different ways
- A standard has evolved that defines what parameters to use
  - Simplifies back annotation
  - Allows for accelerated models (hard-coded)

VITAL models of gates

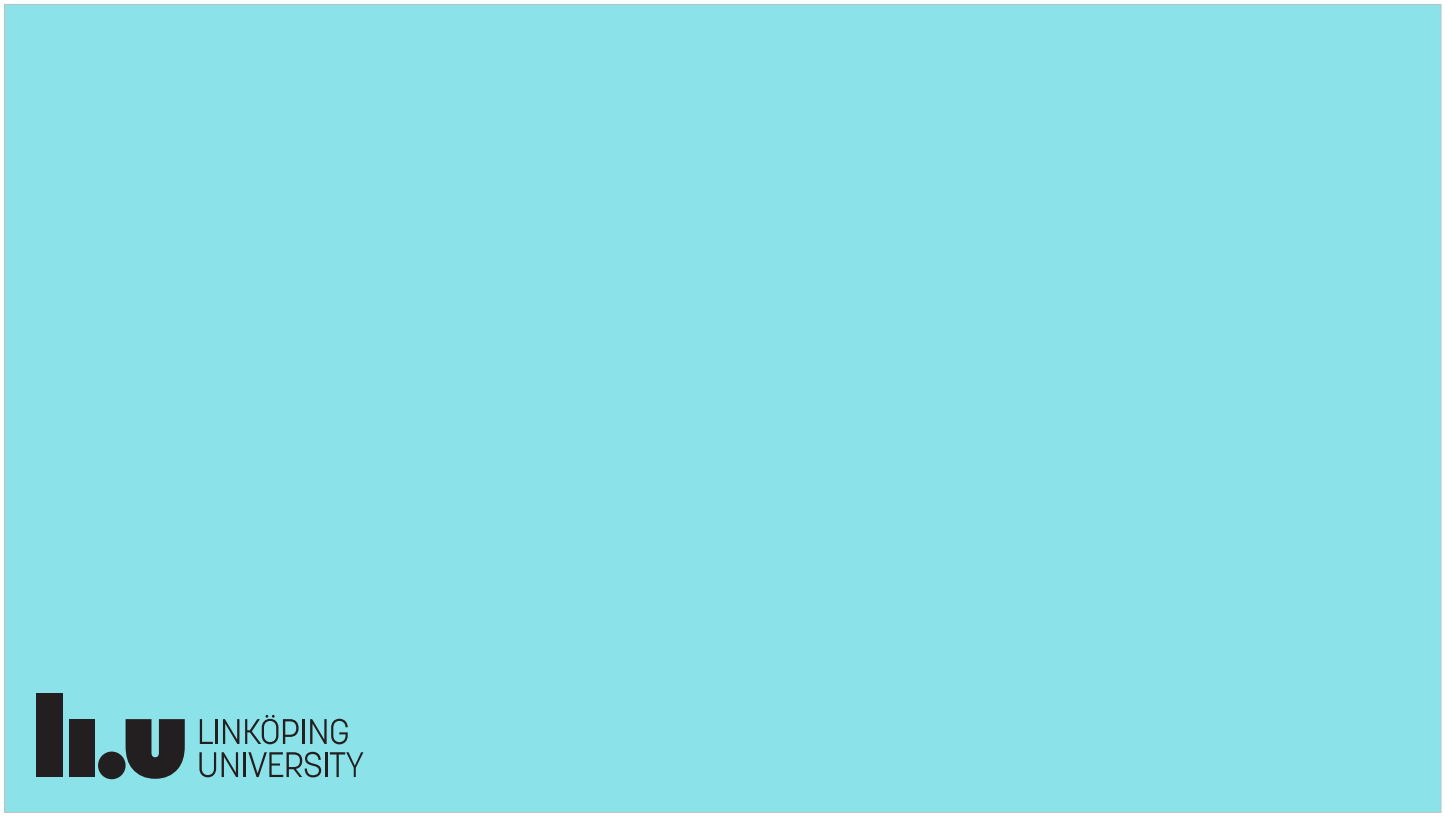
- Three parts: Input delay, Functional and Path delay

```
library IEEE;
use IEEE.VITAL_Primitives.all;
library LIBVUOF;
use LIBVUOF.VTABLES.all;
architecture VITAL of ONAND is
  attribute VITAL_LEVEL1 of VITAL : architecture is TRUE;
  SIGNAL A_ipd  : STD_ULOGIC := 'X';
  SIGNAL B_ipd  : STD_ULOGIC := 'X';
  SIGNAL C_ipd  : STD_ULOGIC := 'X';
  SIGNAL D_ipd  : STD_ULOGIC := 'X';
begin
--  INPUT PATH DELAYS
  WireDelay : block
  begin
    VitalWireDelay (A_ipd, A, tipd_A);
    VitalWireDelay (B_ipd, B, tipd_B);
    VitalWireDelay (C_ipd, C, tipd_C);
    VitalWireDelay (D_ipd, D, tipd_D);
  end block;
```

```
--  BEHAVIOR SECTION
  VITALBehavior : process (A_ipd, B_ipd, C_ipd, D_ipd)
   -- functionality results
    VARIABLE Results: STD_LOGIC_VECTOR(1 to 1):=(others => 'X');
    ALIAS Y_zd : STD_LOGIC is Results(1);
   -- output glitch detection variables
    VARIABLE Y_GlitchData : VitalGlitchDataType;
  begin
   --  Functionality Section
    Y_zd := (NOT ((D_ipd) AND ((B_ipd) OR (A_ipd) OR C_ipd))));
   --  Path Delay Section
    VitalPathDelay01 (
      OutSignal => Y,
      GlitchData => Y_GlitchData,
      OutSignalName => "Y",
      OutTemp => Y_zd,
      Paths => (0 => (A_ipd'last_event, tpd_A_Y, TRUE),
                1 => (B_ipd'last_event, tpd_B_Y, TRUE),
                2 => (C_ipd'last_event, tpd_C_Y, TRUE),
                3 => (D_ipd'last_event, tpd_D_Y, TRUE)),
      Mode => OnDetect,
      Xon => Xon,
      MsgOn => MsgOn,
      MsgSeverity => WARNING);
  end process;
end VITAL;
```


Detection of timing errors

- Input path delay: Transport delay dependent on previous value and wire delay
- Functional part. Boolean expression or lookup tables for fast simulation
- Path delay: output delay, glitch handling
- Models often includes error detection
  - Short spikes, short setup/hold timing etc.
  - Unacceptable values (Z or X)
  - Unacceptable input combinations (both set and reset active on SR flipflop)


LINKÖPING
UNIVERSITY