

TSTE12 Design of Digital Systems Lecture 3

Kent Palmkvist



Agenda

- Practical issues
 - Material
 - Lab 1 requirements
 - Intro to testbench
- Project task intro
- Introduction to VHDL, continued
 - Processes, subprograms etc.
 - Timing

TSTE12 Deadlines Y,D,ED

- Group definitions Thursday 1 September (afternoon)
 - On web, include supervisor assignment
 - Friday 2 September (possible also monday 4/9): First meeting with supervisor
 - Determine project manager (contact person)
 - Questions (short meeting)
 - Tuesday 5 September: First version of requirement specification
 - Hint: Deadline means "no later than", i.e., allowed to complete tasks before these dates
-

TSTE12 Deadlines MELE, erasmus

- Group definition Wednesday 6 September (afternoon)
 - On web, include supervisor assignment
 - Friday 8 September: First meeting with supervisor
 - Determine project manager (contact person)
 - Question (short meeting)
 - Tuesday 12 September: First version of requirement specification
-

Useful resources

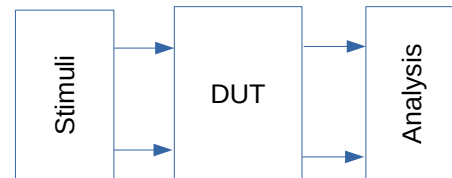
- Lab hardware description (on the linux machines)
 - /courses/TSTE12/material/DE2-115_SystemCD_v3.0.6
 - Lab board datasheets, schematic of hardware

Lab 1 requirements

- Working design (both simulation and hardware)
- Only use `sys_clk` as (the only) clock signal
 - This is the 50 MHz `sys_clk` on the board
- Testbench testing two different key sequences
 - Test after each key scan code input has been generated
 - A text message in the simulator window should be displayed after each scan code test telling the user if test was ok or failed.
- Hint: No need to count the number of received bits, or verify parity and stop bits. Is OK to have intermediate flickering on the display while data is received

VHDL Model Testing (Testbench approach)

- Describe test setup using VHDL
 - Standard approach
 - Test of design will be independent of simulation tool
- Testbench VHDL code contains
 - Stimuli generation
 - Design under test (DUT)
 - Analysis of design output



Programming hardware in the lab

- FPGA can be configured from any of the computers over the network
 - Start Quartus Programmer (program quartus_pgmw) from mentorskal terminal
 - Select hardware setup
 - Select “Add server”
 - Specify machine name to which the board is connected, use the password described on the whiteboard at the door
 - muxen2-009, muxen2-012, muxen2-015
 - Make sure the correct machine and board is defined as hardware

Useful linux commands

- Remember to always work from the “mentorskal” terminal window (except for handins)
- Sets the access rights correctly to support sharing of files within the groups (but not between the groups)

`ls -la file_or_directory`

shows owner and access rights of files

`chmod g+rx filename_or_directory`

let all members of the group have full access to the file or directory (allow others to change files)

Project information

- Project directory will be available for all groups
 - /courses/TSTE12/proj/projgrpXX
 - Access only to members of the group
 - This directory must be used
- Project document templates
 - /courses/TSTE12/material/project/LIPS-templates
- Project functionality
 - Project directive (description) on the web
 - Different priority (mandatory, selective, optional)

Project directive, high priority

- Volume & balance control (atleast 10 levels each)
 - Volume is logarithmic
 - Simplifications are often possible
 - Multiplication is generally expensive
 - Amplification (factor > 1) may lead to overflow!
- A/D and D/A implemented in separate codec chip
 - Supports volume control (not necessary to feed audio into FPGA)
 - Still need to initialize the codec chip (implement I2C communication)

Optional requirement, medium priority

- Each project group make their own choice
 - No problem if more than one group selects the same project
 - Chose one or at most two medium priority requirements
- Oscilloscope
 - Vertical vs horizontal
 - Zoom, color coding, average
 - FFT
 - Complex algorithm
 - Use available designs if possible

Optional requirement, medium priority

- Echo
 - Long echo requires large memory (2MByte SRAM)
 - Number representation
 - 44 KHz, 20 bit, stereo => 220KByte/s
 - Dynamic range, Accuracy
 - 2's complement/Floating point
 - Controls?
 - On/off, Length, Strength

Optional requirement, medium priority

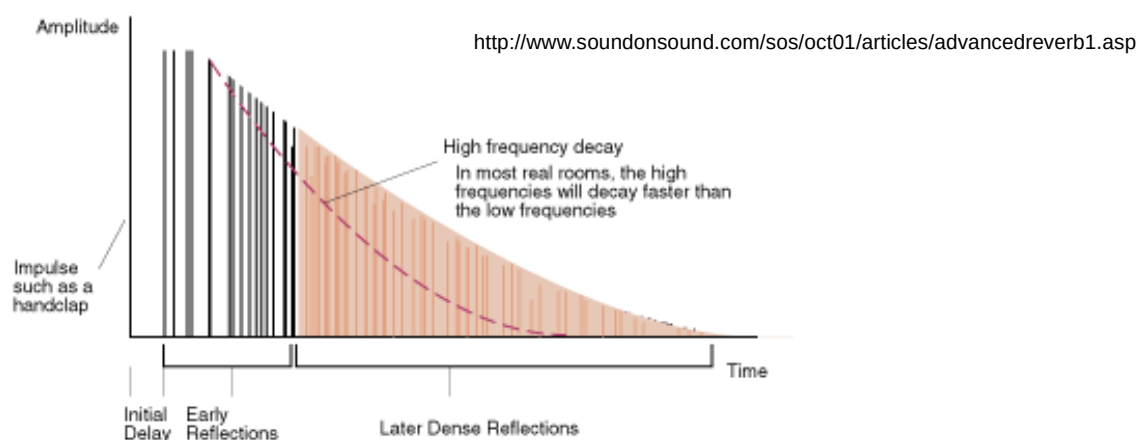
- Signal level indication (logarithmic scale)
 - Peak, average
 - Screen update ~ 50Hz
 - Do not want a flickery output
 - Large number of samples to average on
- Peak indication
 - move slowly towards zero
 - quickly updated when large peak found

Optional requirement, medium priority

- Loudness, supression of mono
 - Loudness control => attenuate medium frequencies
- Equalizer
 - Requires understanding of digital filters
- Test filter algorithms using Matlab
- Other sound modification algorithms possible
 - Reverb, Flange
 - Pitch shifts
 -

Big hall simulation (reverb)

- Principle: Echoes from walls and objects



Reverb, cont.

- Three parts of the sound
 - Direct path
 - Early reflections
 - Decaying mix of reflexions
- Time in the echoes are short
 - Example: 5, 10, 20, 30, 45, 60, 75 ms versions added
 - Create attenuated feedback

Optional requirement, low priority

- Try to think of a lot of optional features
 - Color coding, background image
 - Animations, different VGA screen backgrounds
 - Special keys to use (e.g., arrow keys)
 - Change parameters at run time
 - Additional volume/balance steps
 -
- Try to optimize the size of the design
- Many clock cycles between input samples

Process statement

```
[LABEL:] process (SENSITIVITY_SIGNAL_LIST)
  -- constants, types, subtypes, subprograms
  -- variables (NO signals)
begin
  -- sequential statements
end process;
```

- Process always executed once at simulation start. Then whenever events occur on signals in the sensitivity list
- Variables are static
 - Initialized at simulation start
 - Keep their values between process activations

Process statements, cont.

- Can combine process statements with structural code in the same architecture
- Processes without a sensitivity list will automatically restart at the end of a process
 - Must have some way to stop simulation/wait some time to avoid an infinite loop (simulation appear to be hanged)
- Processes can not both have a sensitivity list and wait statements in the same process

Sequential code

- Processes and subprograms have sequential code
 - One statement after another is executed in order
 - Most similar to "ordinary" computer code
 - Simulation time does not increase while executing the statements
 - Exception is the wait statement

Sequential control statements: wait

- Used in processes and subprograms
- Examples
 - wait on x,y until z=0 for 100 ns;
 - wait until event on x or y while z≠0 or max 100 ns)
 - wait for 100 ns;
 - wait on a,b,c; -- wait for at least one event on a, b or c
 - wait until z=0;
 - wait; -- infinite wait

Sequential control: if

```
if CONDITION1 then
  -- sequence of statements 1
elsif CONDITION2 then
  -- sequence of statements 2
  -- any number of elsif clauses
else
  -- last sequence of statements
end if;
```

- Indentation not important (not like python)
- CONDITION must return boolean (not enough with a bit)

Sequential control: case

```
case EXPRESSION is
  when CHOICE1 => -- sequence of statements1
  when CHOICE2 => -- sequence of statements 2
  when others => -- last sequence of statements
end case;
```

- All possible choices must be covered once
 - Others catch all choices not covered earlier
- Choices may be a list (e.g., when 0|1|2 =>)

Sequential control: loop

```
for NAME in RANGE loop
  -- sequence of operations
end loop;
while CONDITION loop
  -- sequence of operations
end loop;
loop
  -- sequence of statements
end loop;
```

- Controlling loop behavior
 - next [loop_label] [when CONDITION];
 - Skip the rest of the loop body
 - exit [loop_label] [when CONDITION];
 - Terminate the loop

Various statements

Null

- Used to complete syntax requirement, e.g. in case statements when a choice should not do anything.

Subprograms

- Functions and procedures
- Declared in declaration region of architecture, process, block, or other subprograms.
- Variables are dynamic (initialized at every call)
- Functions
 - Always returns a value (must be used in an expression)
 - Never modifies its parameters (all parameters are inputs)
 - No side effects allowed
 - Can not contain wait statements

Functions

```
function FUNCTION_NAME (FORMAL_PARAMETER_DECLARATIONS)
  return RETURN_TYPE is
    -- constant and variable declarations (NO signals)
  begin
    -- sequential statements
    return (RETURN_VALUE);
  end FUNCTION_NAME;
```

- Must always return a defined value

Procedures

```
procedure PROCEDURE_NAME
  (FORMAL_PARAMETER_DECLARATIONS)
  -- Procedure declaration part
  -- constant and variable declarations (NO signals)
begin
  -- sequential statements
end PROCEDURE_NAME;
```

- Formal parameters can be in, out, or inout (default in)
- May contain wait statements (but not if called from a function)
- Procedures can modify its formal parameters (no return value)

Libraries

- All design units are stored in libraries. Accessing a library requires definition.

```
library LIBRARY_NAME;
```

- How this is done in the file system is not defined in VHDL

- Libraries have only logical names

- Simulator and synthesis uses tool-dependent mappings from logical to physical directories (possible including binary format file)

- Allows VHDL source code to be run on different platforms without rewrite

- Two predefined libraries: WORK STD

Packages

- Used in similar way to header files or libraries in other languages
 - Repeat declarations and definitions
- Divided into the declaration and body
 - Declaration part consists of visible declarations when using the package
 - Constants, types, subprograms
 - Declaration parts contains subprogram names and parameter list, but no code
 - Body is visible only within the package. Includes the code for the subprograms
 - Only one possible body definition

Package example

```
package HANDY is
  subtype BITVECT3 is BIT_VECTOR(0 to 2);
  subtype BITVECT2 is BIT_VECTOR(0 to 1);
  function MAJ3(X: BIT_VECTOR(0 to 2) return BIT;

  --- Other declarations -----
end HANDY

package body HANDY is
  function MAJ3(X: BIT_VECTOR(0 to 2))
    return BIT is
  begin
    return (X(0) and X(1)) or (X(0) and X(2)) or (X(1) and X(2));
  end MAJ3;

  --- Other subprogram declarations -----
end HANDY;
```


VHDL cont., Packages

- Packages included in a design by the use clause
 - use `LIBRARY_NAME.PACKAGE_NAME.ELEMENT_NAME;`
 - Element name ALL matches complete body declaration
- Standard packages: `std.standard`
 - boolean, bit, character, severity level, string, bit_vector, time
 - Automatically included
- Packages and libraries used to add new functionality to the language
 - Additional datatypes to support e.g., tristate levels, mathematical functions, etc.

Packages, 1164 standard

- `IEEE.STD_LOGIC_1164`
 - `std_logic` data types, overloaded boolean functions, conversion functions, edge detection functions
- `IEEE.NUMERIC_STD`
 - Overloaded arithmetic functions for `std_logic_vector` and conversion functions to/from integer
 - Defines `UNSIGNED` and `SIGNED` datatypes
 - Can be converted to/from `std_logic_vector` through typecast, e.g.,
`std_logic_vector_signal <= std_logic_vector(unsigned_vectorsignal)`
- `IEEE.NUMERIC_BIT`
 - Allows numeric processing of bitvectors

Configurations

- Defines which entity in a library should be used in the structural code
- Structural example: How to know which gate is used?
 - Component declaration name does not need to exist
 - Assume names are XOR2 and NAND2
- Configuration maps between entity and architecture in library and component name in structure

```
architecture STRUCTURAL of ONES_COUNT is
    component XOR_GATE
        port (X,Y : in bit; O : out bit);
    end component;

    component NAND_GATE
        port (X,Y : in bit; O : out bit);
    end component

    signal I1, I2, I3 : bit;

begin
    U1 : XOR_GATE port map(A(0),A(1),I1);
    U2 : XOR_GATE port map(I1,A(2),C(0));
    U3 : NAND_GATE port map(A(0),A(1),I2);
    U4 : NAND_GATE port map(A(2),I1,I3);
    U5 : NAND_GATE port map(I2,I3,C(1));

end STRUCTURAL;
```

Configuration example

- Add configuration statement to select exact instance and architecture for each instance.
- Drawback: require recompilation if library entity mapping is changed

```
architecture STRUCTURAL of ONES_COUNT is

    component XOR_GATE
        port (X,Y : in bit; O : out bit);
    end component;

    component NAND_GATE
        port (X,Y : in bit; O : out bit);
    end component

    for U1 : XOR_GATE use entity work.xor2(behav);
    for U2 : XOR_GATE use entity work.xor2(behav);
    for U3 : NAND_GATE use entity work.nand2(behav);
    for U4 : NAND_GATE use entity work.nand2(behav);
    for U5 : NAND_GATE use entity work.nand2(behav);

    signal I1, I2, I3 : bit;

begin

    U1 : XOR_GATE port map(A(0),A(1),I1);
    U2 : XOR_GATE port map(I1,A(2),C(0));
    U3 : NAND_GATE port map(A(0),A(1),I2);
    U4 : NAND_GATE port map(A(2),I1,I3);
    U5 : NAND_GATE port map(I2,I3,C(1));

end STRUCTURAL;
```

Configurations, cont.

- Possible to create a configuration as a separate unit (here named STRUCT_BEHAV_XOR)

```
configuration STRUCT_BEHAV_XOR of ONES_COUNT is
  for BEHAV
    for U1:XOR_GATE use entity work.EXORGATE(BEHAV);
  end for;
    for U2:XOR_GATE use entity work.EXORGATE(BEHAV);
  end for;
end STRUCT_BEHAV_XOR;
```

- Multiple configurations possible for the same entity (similar to having multiple architectures)
 - Useful for testbench where both behavioral model and final netlist version of the device should be tested

Visibility

- Same name may be used for multiple purposes
 - Doing this may hide some units depending on the situation

```
package SIG is
  signal X: INTEGER:= 1;
end SIG;

use work.SIG.all;
entity Y is
  signal X: INTEGER:= 2;
end Y;
```

```
architecture Z of Y is
  signal Z1,Z2,Z3,Z4,Z5: INTEGER:= 0;
  function R return INTEGER is
    variable X: INTEGER := 3;
  begin
    return X; -- Returns value of 3.
  end R;
begin
  B: block
    signal X: INTEGER := 4;
    signal Z6: INTEGER := 0;
  begin
    Z6 <= X + Y.X; -- Z6 = 6
  end block B;

  P1: process
    variable X: INTEGER :=5;
  begin
    Z5 <= X; -- Z5=5
    wait;
  end process;

  Z1 <= work.SIG.X; -- Z1=1
  Z2 <= X; -- Z2=2
  Z3 <= R; -- Z3=3
  Z4 <= B.X; -- Z4=4
end Z;
```

Advanced feature: Overloading

- Can change meaning of literals, names of operators, functions, and procedures
- Correct version found by matching both name and parameter types
- If multiple identical declarations => latest visible one is used
- Useful feature for extending the language to support new datatypes or functions
 - E.g. extend to support complex valued real numbers
- IEEE standard 1164 is an example of this
 - Defines new datatypes as well as how comparison, assignments, arithmetic is done using these datatypes

IEEE std_logic_1164 package

- Define a new datatype and support functions
 - the std_logic data type
 - Overloads all common functions used for bits and bitvectors
 - AND, OR, NOT, XOR
 - <, >, = etc.
 - Conversion functions between std_logic and bit
- Netlists generated by synthesis use std_logic for all signals

```

type std_ulogic is ('u', -- uninitialized
                   'x', -- forcing unknown
                   '0', -- forcing 0
                   '1', -- forcing 1
                   'z', -- high impedance
                   'w', -- weak unknown
                   'l', -- weak 0
                   'h', -- weak 1
                   '- ' -- don't care
                   );

```

Overloading example

- Example: Want to support new data type
 - $F \leq (A \text{ and } B) \text{ or } (C \text{ and } D)$;
- new version of and
 - Uses table lookup with non-numeric indexes
 - Index number is based on position in definition enumeration of MVL4

```

type MVL4 is ('X', '0', '1', 'Z');

function "and" (L, R: MVL4) return MVL4 is
-- Declare a two-dimensional table type.
type MVL4_TABLE is array (MVL4,MVL4) of MVL4;
-- truth table for "and" function
constant table_AND: MVL4_TABLE :=
-----
-- | X      0      1      Z
-----
((('X', '0', 'X', 'X'), -- | X |
 ('0', '0', '0', '0'), -- | 0 |
 ('X', '0', '1', 'X'), -- | 1 |
 ('X', '0', 'X', 'X'))); -- | Z |
begin
  return table_AND(L,R);
end "and";

```

File I/O

- Possible to read or write a file (1993 allow both on same file)
- Formatted IO
 - Not generally human readable (platform dependent)
- TEXT IO
 - Human readable
- Special package includes definitions
 - STD.TEXTIO
 - Functions for open file, read a complete line, and read individual data from the line

Later revisions

- Mostly simplifications and additional function support
- 1993:
 - 8-bit ASCII, identifier restrictions relaxed, declarations simplifications
 - Shared variables (global variables outside processes).
 - Improved reporting in assert statements
- 2008:
 - Simplified sensitivity lists (keyword all to include all signals used)
 - Simplified conditions, allow bit and std_logic values as result of condition
 - Read of output ports on entity

VHDL timing and concurrency

- Simulation of concurrent events (hardware) on a sequential computer
- Must have the same result from simulation independent of execution order of individual event
- Delay is an important property of hardware that must be simulated

Signals vs Variables

- Electronic signals can not change values in 0 seconds
 - Always slopes on voltages going from 0 to 1
- Common sequential code assumes variables are updated before next statement is executed
- Expect different result depending on if variables or signals are used
- Both variables and signals can be used in synthesized code

Signal vs Variable example

- Inputs with changing value at different times

X:	1	4	5	5	3
Y:	2	2	2	3	2
Z:	0	3	2	2	2
	initial	t1	t1+2	t1+4	t1+6

- Result depends on if signals or variables as assigned

AS <= X*Y after 2 ns;

BS <= AS+Z after 2 ns;

AS:	2	2	8	10	15
BS:	2	2	5	10	12
	initial	t1	t1+2	t1+4	t1+6

AV := X*Y;

BV := AV + Z;

AV:	2	8	10	15	6
BV:	2	11	12	17	8
	initial	t1	t1+2	t1+4	t1+6

Signal assignment with delta delay

- Minimum delay is a delta delay
- Delta delay is > 0 s but much smaller than the minimum timestep of the simulator

X:	1	4	4	4
Y:	2	2	2	2
Z:	0	3	3	3
initial		t1	t1+delta	t1+2*delta
AS	$\leq X*Y;$			
BS	$\leq AS+Z;$			
AS:	2	2	8	8
BS:	2	2	5	11

Delta delay

- Can not be explicitly specified
- Delta delays will never add up to a simulation delay in seconds (standard time)
- Sometimes referred to as Macro (simulation time) and micro (delta delays) timing.
- Time may stand still in simulation by continuous signal updates
 - Example: process triggered by a signal that it is updating
 - Combinatorial loops without macro delay in assignments
 - Delta delay is increasing but not the simulation time

Simulation models

- Delta delay only
 - Functional verification of models
- Standard time unit delay only
 - Validate system timing
- Mixed
 - Delta delay where delay is not important
 - Standard time unit delay where delay is significant
 - Study system timing